

Congruent triangles in arrangements of lines

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Abstract

We study the maximum number of congruent triangles in finite arrangements of ℓ lines in the Euclidean plane. Denote this number by $f(\ell)$. We show that $f(5) = 5$ and that the construction realizing this maximum is unique, $f(6) = 8$, and $f(7) = 14$. We also discuss for which integers c there exist arrangements on ℓ lines with exactly c congruent triangles. In parallel, we treat the case when the triangles are faces of the plane graph associated to the arrangement (i.e. the interior of the triangle has empty intersection with every line in the arrangement). Lastly, we formulate four conjectures.

Keywords: Arrangement, congruent triangles.

Math. Subj. Class.: 52C10, 52C30

1 Introduction

A problem from mathematical folklore asks for bounding four congruent triangles with six matchsticks. This is easily done, and left to the reader. Quite naturally, one can ask whether more congruent triangles may be formed by using the same six matchsticks. It seems that this particular problem has not been treated in the literature. Our main focus lies on constructing planar arrangements in which a fixed number ℓ of lines bound as many congruent triangles as possible. For an excellent overview on arrangements and spreads, see Grünbaum's [11]. The results presented in this article are complementary to work of Erdős and Purdy [4, 5] on sets of n points—see also [6].

In this paper, everything happens in \mathbb{R}^2 . An *arrangement* (of lines) \mathcal{A} shall be a finite family of ℓ lines L_1, \dots, L_ℓ . In the following, we will ignore the case when there exists a point common to all lines, and thus assume that $\ell \geq 3$. Denote by \mathfrak{A}_ℓ the set of all arrangements of ℓ lines.

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We associate to $\mathcal{A} \in \mathfrak{A}_\ell$ a graph $\Gamma_{\mathcal{A}}$: the vertices of $\Gamma_{\mathcal{A}}$ correspond to the intersection points of lines from \mathcal{A} and the edges of $\Gamma_{\mathcal{A}}$ correspond to the line-segments between these vertices. $\Gamma_{\mathcal{A}}$ is a plane graph. The vertices, edges, and faces of $\Gamma_{\mathcal{A}}$ are also said to belong to \mathcal{A} .

A *triangle* in $\mathcal{A} \in \mathfrak{A}_\ell$ shall be the convex hull of the set of intersection points of three non-concurrent pairwise non-parallel lines in \mathcal{A} . Denote by $F^{\mathcal{A}}$ the set of all triangles in \mathcal{A} . Whenever $\Delta_1, \Delta_2 \in F^{\mathcal{A}}$ are congruent we write $\Delta_1 \sim \Delta_2$. Let $F_1^{\mathcal{A}}, \dots, F_p^{\mathcal{A}}$ be the equivalence classes with respect to \sim such that $|F_1^{\mathcal{A}}| \geq \dots \geq |F_p^{\mathcal{A}}|$. Here, $|M|$ denotes the cardinal number of M . We call a triangle $\Delta \in F^{\mathcal{A}}$ *facial* if it is a face of $\Gamma_{\mathcal{A}}$, i.e. $L \cap \text{int } \Delta = \emptyset$ for all $L \in \mathcal{A}$. Let $G^{\mathcal{A}} \subset F^{\mathcal{A}}$ be the set of all facial triangles in \mathcal{A} , and, as before, let $G_1^{\mathcal{A}}, \dots, G_q^{\mathcal{A}}$ be the equivalence classes with respect to \sim such that $|G_1^{\mathcal{A}}| \geq \dots \geq |G_q^{\mathcal{A}}|$. Put

$$f(\ell) = \max_{\mathcal{A} \in \mathfrak{A}_\ell} |F_1^{\mathcal{A}}| \quad \text{and} \quad g(\ell) = \max_{\mathcal{A} \in \mathfrak{A}_\ell} |G_1^{\mathcal{A}}|.$$

We shall also be considering restrictions relative to a certain arrangement $\mathcal{A} \in \mathfrak{A}_\ell$, namely, for $k \leq \ell$,

$$f_{\mathcal{A}}(k) = \max_{\mathcal{B} \subset \mathcal{A}, \mathcal{B} \in \mathfrak{A}_k} |F_1^{\mathcal{B}}| \quad \text{and} \quad g_{\mathcal{A}}(k) = \max_{\mathcal{B} \subset \mathcal{A}, \mathcal{B} \in \mathfrak{A}_k} |G_1^{\mathcal{B}}|.$$

We call an arrangement $\mathcal{A} \in \mathfrak{A}_\ell$ *f-optimal* (*g-optimal*) if $|F_1^{\mathcal{A}}| = f(\ell)$ ($|G_1^{\mathcal{A}}| = g(\ell)$). If \mathcal{A} is both *f-optimal* and *g-optimal*, we simply write *optimal*. A triangle from $F_1^{\mathcal{A}}$ or $G_1^{\mathcal{A}}$ is said to be *good*. Note that $F_1^{\mathcal{A}}$ and $G_1^{\mathcal{A}}$ need not be unique. In that case, one makes a choice clearly defining $F_1^{\mathcal{A}}$ and $G_1^{\mathcal{A}}$. The edges and angles of a good triangle will be called *good*, too.

An arrangement is *simple* if no three lines are concurrent. The lines of an arrangement are in *general position* if no two lines are parallel and no three lines are concurrent. Two arrangements \mathcal{A} and \mathcal{B} are *combinatorially equivalent* if their associated graphs $\Gamma_{\mathcal{A}}$ and $\Gamma_{\mathcal{B}}$ are isomorphic. (Note that, as mentioned above, we do not consider line arrangements in which all lines meet in a single point.) $\mathcal{A} \in \mathfrak{A}_\ell$ is *c-unique* if there exists no $\mathcal{B} \in \mathfrak{A}_\ell$ such that \mathcal{A} and \mathcal{B} are (a) not combinatorially equivalent and (b) $|H_1^{\mathcal{A}}| = |H_1^{\mathcal{B}}|$, where H is F or G . In the same vein, we say that two arrangements $\mathcal{A} \in \mathfrak{A}_\ell$ and $\mathcal{B} \in \mathfrak{A}_\ell$ are *g-equivalent* if \mathcal{A} can be obtained from \mathcal{B} by translation, rotation, reflection, and scaling. $\mathcal{A} \in \mathfrak{A}_\ell$ is *g-unique* if there exists no $\mathcal{B} \in \mathfrak{A}_\ell$ such that \mathcal{A} and \mathcal{B} are (a) not g-equivalent and (b) $|H_1^{\mathcal{A}}| = |H_1^{\mathcal{B}}|$, where H is F or G . A few examples: Three lines in general position yield an arrangement that is c-unique, but not g-unique; any arrangement on four lines that forms exactly two congruent triangles is not c-unique (and thus cannot be g-unique); finally, as we shall see in Theorem 3.5, the *f-optimal* arrangement from Figure 2 (b) is g-unique (and thus c-unique).

$F(\ell)$ ($G(\ell)$) is defined as the set of all integers u such that there exists an arrangement on ℓ lines having exactly u congruent triangles (congruent facial triangles). We write $[s..t]$ for the set of all integers u with $s \leq u \leq t$, put H for F or G , and

$$h = \begin{cases} f & \text{if } H = F, \\ g & \text{if } H = G. \end{cases}$$

Whenever $H(\ell) = [0..h(\ell)]$, we say that $H(\ell)$ is *complete*. In the following, we will tacitly use the fact that $G(\ell) \subset F(\ell)$. We call an arrangement $\mathcal{A} \in \mathfrak{A}_\ell$ *1-extendable* if there exists a line L such that $|H_1^{\mathcal{A} \cup L}| = |H_1^{\mathcal{A}}| + 1$.

2 Preparation

We briefly concern ourselves with the following question, since it will shorten later arguments. What if we drop the condition that the triangles need to be congruent? Kobon Fujimura asked in 1978 in his book “The Tokyo Puzzle” [8]—see also [10, pp. 170–171 and 178]—what the maximum number $K(\ell)$ of facial (not necessarily congruent!) triangles realisable by ℓ lines in the plane is. (Grünbaum treated this problem before Fujimura, but he might have only been interested in arrangements in the projective plane [6].) Recently, Bader and Clément [1], improving upon a result of Tamura, showed the following.

Lemma 2.1 (Bader and Clément).

$$K(\ell) \leq \left\lfloor \frac{\ell(\ell - 2)}{3} \right\rfloor - \mathbf{I}_{\{\ell : (\ell \bmod 6) \in \{0, 2\}\}}(\ell),$$

where \mathbf{I} denotes the indicator function.

Many arrangements have been constructed in order to find solutions to Fujimura’s problem. Fujimura himself gave an example which shows that $K(7) \geq 11$, although it was thought for many years that $K(7) = 10$. In 1996, Grabarchuk and Kabanovitch [13] gave two 10-line, 25-triangle constructions, whereas Lemma 2.1 gives $K(10) \leq 26$. Whether $K(10)$ is 25 or 26 is unknown. Other 10-line, 25-triangle arrangements were found independently by Grünbaum [12, p. 400], Wajnberg, and Honma (see [15] for more details). Good overviews of the best (i.e. the greatest number of triangles for a fixed number of lines) known arrangements can be found in [14] and [15].

Table 1: Bounding $K(\ell)$ for $\ell \leq 12$.

ℓ	3	4	5	6	7	8	9	10	11	12
Bader-Clément bound	1	2	5	7	11	15	21	26	33	39
Best known arrangement	1	2	5	7	11	15	21	25	32	38

Füredi and Palásti [9] construct an arrangement proving $K(\ell) \geq \ell(\ell - 3)/3$. See also the article of Forge and Ramírez Alfonsín [7].

We continue with a series of lemmas. Lemma 2.2 is stated without its straightforward proof, but we present the heptagonal case in Figure 1 and its caption.

Lemma 2.2. *The ℓ lines bounding a regular ℓ -gon determine exactly 2ℓ congruent triangles if $\ell \geq 7$, and therefore $2\ell \in F(\ell)$ and $f(\ell) \geq 2\ell$. With the same construction we obtain for $\ell \geq 5$ that $\ell \in G(\ell)$ and $g(\ell) \geq \ell$.*

Lemma 2.3. *Let $\mathcal{A} \in \mathfrak{A}_\ell$ and $h \in \{f, g\}$. Then, for $3 \leq k \leq \ell - 1$,*

$$h_{\mathcal{A}}(\ell) \leq \frac{\ell(\ell - 1)(\ell - 2)}{k(k - 1)(k - 2)} \cdot h_{\mathcal{A}}(k).$$

Proof. We observe that every subset of $k < \ell$ lines within the arrangement \mathcal{A} cannot have more than $h_{\mathcal{A}}(k)$ good triangles (good in \mathcal{A}). Counting all together, there are at most

$\binom{\ell}{k} h_{\mathcal{A}}(k)$ good triangles, each appearing several times. Since each of them lies in $\binom{\ell-3}{k-3}$ sub-arrangements of k lines, we obtain

$$h_{\mathcal{A}}(\ell) \leq \frac{\binom{\ell}{k} h_{\mathcal{A}}(k)}{\binom{\ell-3}{k-3}}.$$

Thus, for each k we obtain an upper bound for $h_{\mathcal{A}}(\ell)$. □

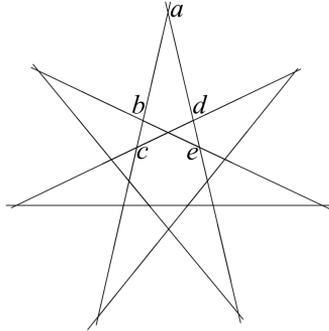


Figure 1: Seven lines bounding a regular heptagon. This arrangement contains fourteen congruent triangles: abe , acd , and their symmetric counterparts obtained by rotating around the barycentre of the heptagon by $2\pi k/7$ for $k = 1, \dots, 6$. This arrangement proves that $f(7) \geq 14$.

Lemma 2.4 is a direct consequence of Lemma 2.3.

Lemma 2.4. *Let $h \in \{f, g\}$. Then*

$$h(\ell) \leq \min_{3 \leq k \leq \ell-1} \frac{\ell(\ell-1)(\ell-2)}{k(k-1)(k-2)} \cdot h(k).$$

3 Results

3.1 Bounds for the general case

Proposition 3.1. *Let $h \in \{f, g\}$, consider $\mathcal{A} \in \mathfrak{A}_\ell$ and $\mathcal{B} \in \mathfrak{A}_k$, and assume that a good triangle of \mathcal{A} is similar to a good triangle of \mathcal{B} . Furthermore, \mathcal{A} and \mathcal{B} each contain two lines intersecting in the boundary of the respective convex hulls of $V(\Gamma_{\mathcal{A}})$ and $V(\Gamma_{\mathcal{B}})$ and forming the same good angle. Then $h(\ell + k - 2) \geq h(\ell) + h(k)$.*

Proof. We scale \mathcal{B} to \mathcal{B}' such that the good triangles of \mathcal{B}' are congruent with the good triangles of \mathcal{A} . Consider the convex hull $C_{\mathcal{A}}$ ($C_{\mathcal{B}'}$) of the intersection points of \mathcal{A} (\mathcal{B}'). Let $p_{\mathcal{A}}$ ($p_{\mathcal{B}'}$) be an intersection point of \mathcal{A} (\mathcal{B}') lying on the boundary of the convex polygon $C_{\mathcal{A}}$ ($C_{\mathcal{B}'}$) and incident with a good angle $\alpha_{\mathcal{A}}$ ($\alpha_{\mathcal{B}'}$) of \mathcal{A} (\mathcal{B}') such that $\alpha_{\mathcal{A}} = \alpha_{\mathcal{B}'}$. Denote with $L_1^{\mathcal{A}}$ and $L_2^{\mathcal{A}}$ ($L_1^{\mathcal{B}'}$ and $L_2^{\mathcal{B}'}$) two of the lines of \mathcal{A} (\mathcal{B}') intersecting at $p_{\mathcal{A}}$ ($p_{\mathcal{B}'}$) and forming the angle $\alpha_{\mathcal{A}}$ ($\alpha_{\mathcal{B}'}$). We can now identify $L_1^{\mathcal{A}}$ with $L_1^{\mathcal{B}'}$ and $L_2^{\mathcal{A}}$ with $L_2^{\mathcal{B}'}$ such that an arrangement \mathcal{C} is obtained in which, seeing \mathcal{A} and \mathcal{B}' as sub-arrangements of \mathcal{C} , no good triangle lies in both \mathcal{A} and \mathcal{B}' . (Note that the number of good triangles in \mathcal{C} may be larger

than the sum of the number of good triangles in \mathcal{A} and \mathcal{B}' , see e.g. Figure 8, in which the arrangements from Figures 1 and 5 (b) are joined: the original arrangements have 14 and 8 good triangles, respectively, but the new arrangement has 26.) \square

We write $\tilde{f}(\ell)$ if we consider the values of $f(\ell)$ only for arrangements whose good triangles are not right triangles.

Proposition 3.2. $\tilde{f}(\ell + 1) \leq \tilde{f}(\ell) + 3(\ell - 1)$ and $f(\ell + 1) \leq f(\ell) + 4(\ell - 1)$.

Proof. Let L be a line being added to $\mathcal{A} \in \mathfrak{A}_\ell$. If \mathcal{A} has the property that $|F_1^{\mathcal{A}}| = |F_2^{\mathcal{A}}|$, then consider henceforth only the triangles in $F_1^{\mathcal{A}}$, as well as their edges, to be good. We denote the lengths of good edges with a, b , and c .

There are at most $\ell - 1$ triangles with an edge of length a on L : there are at most $\ell/2$ lines of one of the two lines needed to make a good triangle with an edge on L , each of these lines is part of at most two triangles with an edge of length a on L , and it is impossible for there to be exactly $\ell/2$ of them each of which is part of exactly two triangles. Since this can be applied analogously for edges of length b and c , we have $\tilde{f}(\ell + 1) = \tilde{f}(\ell) + 3(\ell - 1)$.

For good triangles that are right triangles, we argue in the same manner and obtain that for each of the three types of good edge (i.e. of length a, b or c) there are at most $4(\ell - 1)/3$ triangles with a good edge of that type on L . \square

Proposition 3.3. $\tilde{f}(\ell) \leq \ell(\ell - 1)$ and $f(\ell) \leq 4\ell(\ell - 1)/3$.

Proof. The idea is the same as the one used in the proof of Proposition 3.2. In the case of non-right triangles, we have established that on each line in \mathcal{A} there are at most $3(\ell - 1)$ good edges. By multiplying this with ℓ , we obtain an upper bound for the number of good edges in \mathcal{A} . Now we divide by three (as there are three edges to each triangle) and have the desired bound. The case of right triangles is settled analogously. \square

All angles in $\mathcal{A} \in \mathfrak{A}_\ell$ equal to one of the angles of a good triangle which is not a right triangle will be called *non-right angles*.

Proposition 3.4. $\tilde{f}(\ell) \leq 2\ell(\ell - 2)/3$ for simple arrangements.

Proof. Consider a simple arrangement \mathcal{A} on ℓ lines admitting a good triangle which is not a right triangle. Let $V(\Gamma_{\mathcal{A}}) = V$, and write V_k for the set of vertices of degree k . As no three lines are concurrent, in $\Gamma_{\mathcal{A}}$ there exist only vertices of degree 2, 3, or 4. Trivially, around a vertex of degree 2 at most one non-right angle resides. Around a vertex of degree 3 likewise (as $\pi/2$ is a right angle, and the sum of two non-right angles must be strictly smaller than π), and around a vertex of degree 4 there may be at most two non-right angles.

Thus, in \mathcal{A} , we have as an upper bound for the maximum number of non-right angles $|V_2| + |V_3| + 2 \cdot |V_4| = |V| + |V_4|$. We have $|V| \leq \ell(\ell - 1)/2$. Also $|V_4| \leq |V| - \ell$, because on every line the first and the last vertex belong to $V_2 \cup V_3$, but any such vertex may appear as first or last vertex on two lines of \mathcal{A} . Thus, the bound is $\ell^2 - 2\ell$. For odd $\ell \geq 5$, this bound is best possible: for the ℓ lines bounding a regular ℓ -gon we have $|V_2| = \ell$, $|V_3| = 0$, and $|V_4| = \ell(\ell - 3)/2$. One non-right angle cannot lie in more than two triangles which are not right triangles, and every triangle requires three angles, whence, the final bound. (In fact, no good angle can lie in more than two good triangles.) \square

3.2 $\ell \leq 5$

We have $f(3) = g(3) = 1$ and $f(4) = g(4) = 2$, and $F(3)$, $G(3)$, $F(4)$, and $G(4)$ are complete. We leave the easy proofs to the reader, but mention that in the 4-line case there exist exactly three combinatorially different solutions with two congruent triangles (these coincide for the g -optimal and the f -optimal case): one with three concurrent lines, one with two parallel lines, and one in general position.

We now focus on the first interesting case: $\ell = 5$.

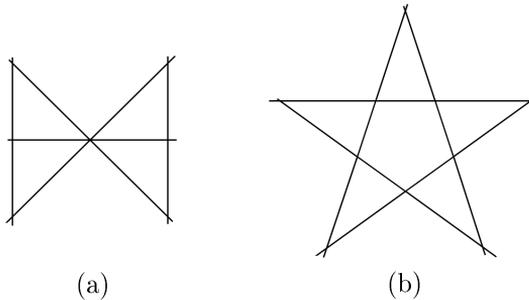


Figure 2: (a) shows a 5-line arrangement with four congruent triangles constructed as follows. Two lines L_1, L_2 orthogonal to a third line L_3 are considered. Let the intersection points be p_1 and p_2 , resp. A fourth and fifth line are considered such that their intersection point is the midpoint of the line-segment p_1p_2 and the angle each forms with L_3 is $\pi/4$. (b) depicts the five lines bounding a regular pentagon. This arrangement contains ten triangles, distributed among two congruence classes of size 5 each.

Theorem 3.5. (i) We have $f(5) = g(5) = 5$ while $F(5)$ and $G(5)$ are complete. Furthermore, the arrangement from Figure 2 (b) is (ii) optimal, and (iii) g -unique among f -optimal and g -optimal 5-line arrangements.

Proof. Figure 2 (b) shows that $g(5) \geq 5$ (whence, $f(5) \geq 5$), with which Lemma 2.1 implies $g(5) = 5$. $f(5) = 5$ follows from Lemma 2.4 (with $k = 4$). We now discuss $G(5)$. We have $G(4) = [0..2] \subset G(5)$. Consider the four lines bounding a square and add the two lines containing the square’s diagonals. By removing one of the four original lines, we have shown that $3 \in G(5)$. Together with the arrangements from Figure 2, we have that $G(5)$ is complete since $g(5) = 5$. Thus, (i) is proven. (ii) follows directly from (i).

We now prove (iii). First, we show that the arrangement from Figure 2 (b) is c -unique. We use the database provided in Christ’s Dissertation [3, Chapter 3.2.5]. (A visualisation of Christ’s results is available in [2]. Note that this does not coincide with Grünbaum’s isomorphism types of arrangements given in [11, p. 5], since Grünbaum discusses the issue in the projective plane, while here we treat the situation in the Euclidean plane.) Among arrangements of five lines in general position, there are exactly six combinatorially different ones, shown in Figure 3. The arrangement in Figure 2 (b) belongs to the combinatorial class (A).

Only the arrangements in (A) contain five facial triangles, i.e. triangles which are faces in the associated graph. We leave to the reader the straightforward proof that among arrangements with five lines not in general position (i.e. containing two parallel lines or three

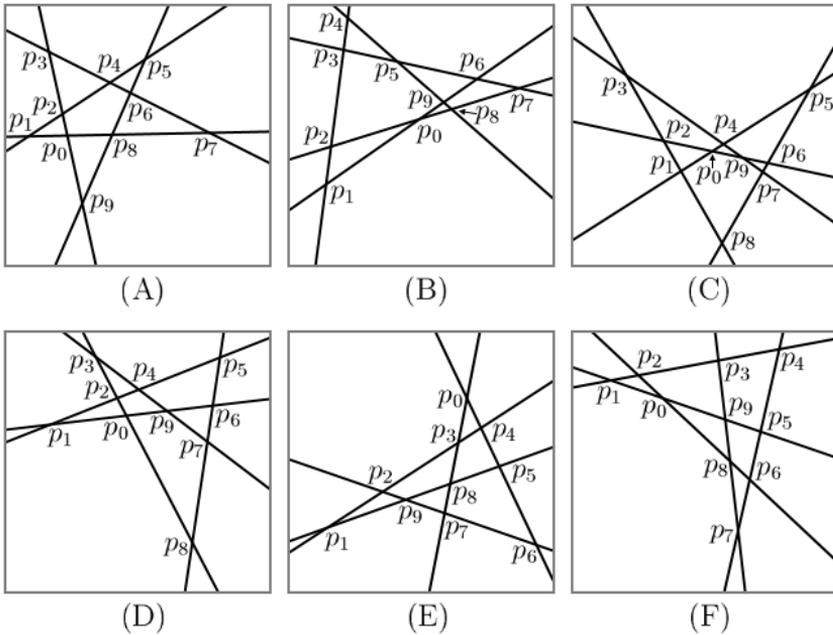


Figure 3: Representations of the six combinatorially different arrangements of five lines in general position.

concurrent lines), there is none featuring five facial triangles. Note that the occurrence of more triangles is impossible due to Lemma 2.1.

We now turn to the case in which triangles are not facial. Let us denote a line-segment between two points x, y with xy and its length with $L(xy)$. We will use the following.

Remark. The sum of the measures α and β of two good angles is π if and only if $\alpha = \beta = \pi/2$.

We write Δ_{ijk} for the triangle with vertices p_i, p_j, p_k . We will tacitly make use of the fact that if in a given arrangement a triangle Δ is strictly contained in a triangle Δ' , then $\Delta \approx \Delta'$ and so Δ and Δ' cannot lie in the same congruence class.

(B) We have $\Delta_{012} \subset \Delta_{149} \cap \Delta_{136}$ (so $\Delta_{149} \approx \Delta_{012} \approx \Delta_{136}$), $\Delta_{345} \subset \Delta_{149} \cap \Delta_{248}$, $\Delta_{569} \subset \Delta_{136} \cap \Delta_{237} \cap \Delta_{578}$, $\Delta_{578} \subset \Delta_{237}$, $\Delta_{067} \subset \Delta_{237}$, $\Delta_{089} \subset \Delta_{248} \cap \Delta_{237} \cap \Delta_{067}$. Due to these inclusion relations, only the following set of triangles may form a congruence class of size 5: $\{\Delta_{149}, \Delta_{136}, \Delta_{248}, \Delta_{578}, \Delta_{067}\}$. Assume it is indeed a congruence class. Thus, all angles around p_6 are right. We apply the Remark to the angles surrounding p_6 . Combining this with $\Delta_{136} \sim \Delta_{067}$ and $p_0p_6 \subsetneq p_1p_6$, we have $L(p_6p_7) = L(p_1p_6)$, $L(p_0p_6) = L(p_3p_6)$, and $L(p_0p_7) = L(p_1p_3)$. p_1p_3 is the hypotenuse in Δ_{136} , but as p_1p_4 is an edge of Δ_{149} and $\Delta_{149} \sim \Delta_{136}$ we obtain a contradiction, since $L(p_1p_4) > L(p_1p_3)$.

(C) We have $\Delta_{012} \subset \Delta_{134} \cap \Delta_{268} \cap \Delta_{378}$, $\Delta_{134} \cup \Delta_{239} \subset \Delta_{378}$, $\Delta_{457} \cup \Delta_{056} \subset \Delta_{158}$, $\Delta_{049} \subset \Delta_{158} \cap \Delta_{239} \cap \Delta_{378} \cap \Delta_{056}$, $\Delta_{679} \subset \Delta_{158} \cap \Delta_{268} \cap \Delta_{457}$. There is no congruence class of size 5.

(D) We have $\Delta_{012} \subset \Delta_{149} \subset \Delta_{156}$, $\Delta_{234} \subset \Delta_{039} \subset \Delta_{378}$, $\Delta_{457} \cup \Delta_{068} \subset \Delta_{258}$, $\Delta_{679} \subset \Delta_{258} \cap \Delta_{457} \cap \Delta_{068}$. Once more, all congruence classes have size at most 4.

(E) We have $\Delta_{129} \subset \Delta_{138} \subset \Delta_{145}$, $\Delta_{237} \subset \Delta_{246}$, $\Delta_{034} \subset \Delta_{058} \subset \Delta_{067}$, $\Delta_{789} \subset \Delta_{237} \cap \Delta_{246} \cap \Delta_{569}$. As above.

(F) We have $\Delta_{012} \subset \Delta_{139} \subset \Delta_{145}$, $\Delta_{238} \cup \Delta_{056} \subset \Delta_{246}$, $\Delta_{678} \subset \Delta_{579} \subset \Delta_{347}$, $\Delta_{089} \subset \Delta_{238} \cap \Delta_{246} \cap \Delta_{056}$. As above.

Let us show that in a 5-line arrangement \mathcal{A} containing two parallel lines or three concurrent lines, no more than four congruent triangles can be achieved. We first assume that \mathcal{A} contains parallel lines L_1, L_2 . If there exists a line L_3 parallel to L_1 , we are done, as in \mathcal{A} there are only at most three triples of lines forming triangles. Thus, w.l.o.g. we are in the situation that a third line, L_3 , intersects L_1 and L_2 . Now assume that a fourth line, L_4 , is parallel to L_3 . Note that L_1, L_2, L_3, L_4 bound zero triangles. In this situation, a fifth line generates at most four new triangles. Thus, L_4 cannot be parallel to L_3 . We have proven that a 5-line arrangement containing three parallel lines or two parallel pairs of parallel lines cannot have more than four congruent triangles.

Denote the open strip bounded by L_1 and L_2 with S , and the complement of its closure by T . Also, let T_1 and T_2 be the connected components of T . In the light of above paragraph, there are three cases (see Figure 4): either (a) L_4 is concurrent with L_1 and L_3 in a point x lying in the closure of T_1 , (b) L_4 intersects L_3 in S , or (c) L_4 intersects L_3 in T_2 . Denote with L_5 the fifth line of \mathcal{A} . We know that L_5 is not parallel to any of the existing four lines. We write Δ_{ijk} for the triangle bounded by the lines L_i, L_j, L_k .

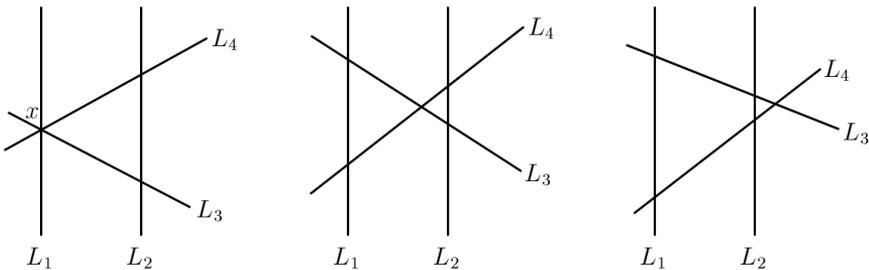


Figure 4: Cases (a)–(c) occurring in the proof of Theorem 3.5.

(a): If $x \in L_5$, then the five lines would bound only three triangles, so we can assume $x \notin L_5$.

Case 1: L_5 intersects both L_3 and L_4 in S . We have six triangles, but $\Delta_{345} = \Delta_{234} \cap \Delta_{145}$ and $\Delta_{245} \subset \Delta_{235}$, so the maximum number of congruent triangles is three.

Case 2: L_5 intersects L_4 in S and L_5 intersects L_3 in $T_1 \cup T_2$. Six triangles appear. Subcase 2.1: L_3 and L_5 intersect in T_1 . But then we have $\Delta_{135} \subset \Delta_{345} \subset \Delta_{235}$. Subcase 2.2: L_3 and L_5 intersect in T_2 . Here, $\Delta_{235} \subset \Delta_{345} \subset \Delta_{135}$, so once more five congruent triangles cannot occur. (If L_5 intersects L_3 in S and L_4 in T_2 , then we are, combinatorially, in the situation treated in Subcase 2.2.)

Case 3: L_5 is concurrent with L_2 and L_3 . Subcase 3.1: All intersection points lie in the closure of S . Five triangles appear, but among them one is a subset of another. Subcase 3.2: L_4 and L_5 intersect in T_2 . We apply the same argument as before. Subcase 3.3: L_4 and L_5 intersect in T_1 . Once more five triangles occur, but one is contained in another.

Case 4: L_5 intersects L_3 and L_4 in points p and p' , respectively, which do not lie in S (since this was covered in Cases 1 and 2). Subcase 4.1: If p and p' lie in T_1 , six triangles appear, but $\Delta_{345} \subset \Delta_{135} \subset \Delta_{235}$. Subcase 4.2: If p and p' lie in T_2 , again six triangles occur, but $\Delta_{245} \subset \Delta_{235} \subset \Delta_{135}$. Subcase 4.3: If $p \in T_1$ and $p' \in T_2$, six triangles are present in the arrangement, but $\Delta_{245} \subset \Delta_{145} \subset \Delta_{345}$, so at most four triangles are congruent.

Case 5: L_5 is concurrent with L_2 and L_4 . Subcase 5.1: All intersection points lie in the closure of S . Subcase 5.1 coincides with Subcase 3.1. Subcase 5.2: L_3 and L_5 intersect in T_2 . But then we are in the same situation as Subcase 3.2.

(b) Let L_3 and L_4 intersect in y . We know that L_5 is not parallel to L_1 . If $y \in L_5$, we obtain six triangles. However, either $\Delta_{135} \cup \Delta_{145} = \Delta_{134}$ and symmetrically $\Delta_{235} \cup \Delta_{245} = \Delta_{234}$ or $\Delta_{134} \cup \Delta_{145} = \Delta_{135}$ and $\Delta_{245} \cup \Delta_{234} = \Delta_{235}$. In either case, the largest congruence class has cardinality at most four. We have treated the cases when L_5 is concurrent with L_1 and L_3 , L_1 and L_4 , L_2 and L_3 , or L_2 and L_4 in (i). We split the remaining cases into four cases according to where the intersection points of L_5 with L_3 and L_4 lie. In each situation, inclusions are given which make the occurrence of a congruence class of cardinality at least five impossible.

Case 1: Both intersection points lie in S . However, we then have $\Delta_{135} \subset \Delta_{145}$, $\Delta_{345} = \Delta_{134} \cap \Delta_{235}$, and $\Delta_{234} \subset \Delta_{245}$.

Case 2: The intersection points of L_5 with L_3 and L_4 lie in T_1 and T_2 , respectively. Then $\Delta_{135} \cup \Delta_{245} \subset \Delta_{345}$, $\Delta_{135} \cup \Delta_{234} \subset \Delta_{235}$, and $\Delta_{134} \cup \Delta_{245} \subset \Delta_{145}$.

Case 3: The intersection points of L_5 with L_3 and L_4 lie in T_1 and S , respectively. We have $\Delta_{135} \subset \Delta_{345} \subset \Delta_{235}$, $\Delta_{245} \subset \Delta_{234}$, and $\Delta_{134} \subset \Delta_{145}$.

Case 4: Both intersection points lie in T_1 . Then $\Delta_{234} \subset \Delta_{235}$, $\Delta_{135} = \Delta_{235} \cap \Delta_{145}$, and $\Delta_{134} \subset \Delta_{145} \subset \Delta_{245}$.

As situation (c) uses very similar arguments, we skip it.

We have shown that no two lines in \mathcal{A} are parallel. Assume now that three lines L_1, L_2, L_3 of \mathcal{A} intersect at a point q . If L_4 contains q as well, the largest congruence class which may be formed by a fifth line has size 2. So $q \notin L_4$ and L_4 is not parallel to any of L_1, L_2, L_3 . W.l.o.g. let the intersection point of L_4 with L_2 lie between the intersection point of L_4 with L_1 and the intersection point of L_4 with L_3 . If there are two coincidences (of three lines)—it is easy to see that there cannot be more—we have three combinatorially different cases. W.l.o.g., in each of them L_1, L_4 , and L_5 shall be concurrent. We denote this intersection point with a , and the intersection point of L_5 with L_2 and L_3 with b and c , respectively. We differentiate the three cases by the order in which the intersection points occur on L_5 .

Case 1: $a - c - b$ (or equivalently $b - c - a$): Eight triangles occur. However, we have $\Delta_{124} \cup \Delta_{234} = \Delta_{134} \subset \Delta_{345}$ and $\Delta_{135} \cup \Delta_{235} = \Delta_{125} \subset \Delta_{245}$. Thus, no five triangles can be congruent.

Case 2: $b - a - c$ (or equivalently $c - a - b$): Again, eight triangles appear, but $\Delta_{245} \cup \Delta_{124} = \Delta_{125}$, $\Delta_{125} \cup \Delta_{135} = \Delta_{235}$, $\Delta_{234} \cup \Delta_{124} = \Delta_{134}$, and $\Delta_{134} \cup \Delta_{135} = \Delta_{345}$.

Case 3: $a - b - c$ (or equivalently $c - b - a$): $\Delta_{124} \cap \Delta_{135} = \Delta_{125}$, $\Delta_{235} \subset \Delta_{234}$, and $\Delta_{245} \subset \Delta_{345}$. Furthermore, every triangle is contained in Δ_{134} .

We are left with the case that there is exactly one coincidence of three lines (namely in q). Once more, several cases occur. We leave them to the reader—treating them is a straightforward task in exactly the same spirit as above paragraphs.

Finally, we prove that the construction from Figure 2 (b) is indeed g-unique. Consider five lines bounding a pentagon P such that we obtain an arrangement \mathcal{A} in combinatorial class (A). This implies that no two lines in \mathcal{A} are parallel. Combinatorially, there are two types of triangles in P : those sharing exactly two vertices (and thus an edge) with P , and those sharing exactly one vertex with P . Due to straightforward inclusion arguments, all triangles in a congruence class of size 5 are of the same type.

Consider the first type, and let Δ be one of these five congruent facial triangles. Denote the angles of Δ incident with a vertex of P with α and β . Applying successively the fact that no two lines in \mathcal{A} are parallel, we obtain that $\alpha = \beta$, so Δ is isosceles. This implies that all angles of P must be equal, and since P is a pentagon, the angles of P measure $3\pi/5$ each. Thus $\alpha = 2\pi/5$ —in particular, Δ is not equilateral. Hence, the sides of P must have equal length, so P is a regular pentagon.

We treat the second case. We see each triangle of the second type as the union of three faces (of $\Gamma_{\mathcal{A}}$): the pentagon P , which lies in all five triangles, and two facial triangles. Since certain pairs of triangles of the second type share a facial triangle, there are at most two congruence classes C_1 and C_2 of facial triangles. Assume $C_1 \neq C_2$. Thus, there exists a triangle Δ of second type containing a facial triangle in C_1 and a facial triangle in C_2 . By considering all five congruent triangles of second type, a contradiction is obtained, since necessarily one of these triangles will contain only triangles from either C_1 or C_2 and thus, it cannot be congruent to Δ . We have proven that all facial triangles are congruent. Now we may argue as in the preceding paragraph. □

3.3 $\ell = 6$

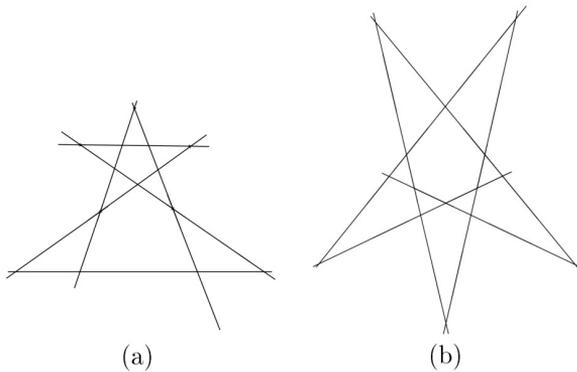


Figure 5: (a) This arrangement is due to Tudor Zamfirescu and shows that $7 \in F(6)$. To the five lines bounding a regular pentagon a sixth line is added which is parallel to one of the five lines such that seven congruent triangles are present. (b) This arrangement proves that $8 \in F(6)$ and $f(6) \geq 8$. In Theorem 3.6 we show that in fact $f(6) = 8$. The arrangement is obtained by considering six of the seven lines bounding a regular heptagon.

Theorem 3.6. *We have $f(6) = 8$, $6 \leq g(6) \leq 7$, $F(6)$ is complete, and $[0..6] \subset G(6)$.*

Proof. The arrangement from Figure 5 (b) proves that $f(6) \geq 8$. Theorem 3.5 (iii) states that there is exactly one f -optimal arrangement on five lines, shown in Figure 2 (b). We

call this arrangement \mathcal{P} . We now show that one cannot produce an arrangement on six lines which has \mathcal{P} as a sub-arrangement and features eight (or more) congruent triangles. Assume there exists such an arrangement \mathcal{A} . Denote the lines of \mathcal{P} by L_1, \dots, L_5 , and the line added to \mathcal{P} in order to obtain \mathcal{A} by L .

First we prove that the addition of L cannot create a “new” congruence class (i.e. a class the triangles of which are non-congruent to every triangle present in \mathcal{P}) of congruent triangles of cardinality at least 8. At least one of the angles $\pi/5, 2\pi/5, 3\pi/5, 4\pi/5$ is good in both \mathcal{P} and \mathcal{A} , since every triangle bounded by L has at least one angle in \mathcal{P} . Among all angles in \mathcal{A} , each of the aforementioned four angles appears at least ten times in five pairs of opposite angles, since \mathcal{P} is a sub-arrangement of \mathcal{A} . Thus, L forms at least three copies of the angle α with the lines L_1, \dots, L_5 , where $\alpha \in \{\pi/5, 2\pi/5, 3\pi/5, 4\pi/5\}$ is fixed. But since the L_i 's are pairwise non-parallel, this is only possible if L is parallel to some L_i . But then the addition of L to \mathcal{P} yields at most six new triangles—too few.

Take the two congruence classes $F_1^{\mathcal{P}}$ and $F_2^{\mathcal{P}}$ such that the triangles in $F_1^{\mathcal{P}}$ are facial in \mathcal{P} , and notice that $F^{\mathcal{P}} = F_1^{\mathcal{P}} \cup F_2^{\mathcal{P}}$ and $|F_1^{\mathcal{P}}| = |F_2^{\mathcal{P}}| = 5$. Thus, L must add at least three triangles to $F_1^{\mathcal{P}}$ or $F_2^{\mathcal{P}}$. As $\pi/5$ belongs to triangles in $F_1^{\mathcal{P}}$ as well as triangles in $F_2^{\mathcal{P}}$, $\pi/5$ is a good angle in \mathcal{A} , so L makes this angle with a line of \mathcal{P} , whence, L is parallel to some L_i , say L_1 . Among all possible positions of L , only three provide new triangles congruent either to a good triangle in $F_1^{\mathcal{P}}$ or to a good triangle in $F_2^{\mathcal{P}}$, see Figure 6. The number of those new triangles is 1, 1, 2, respectively.

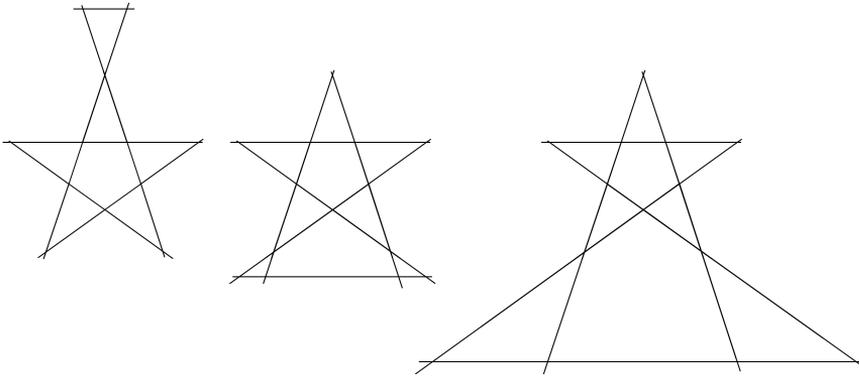


Figure 6: The three essentially different arrangements of five lines bounding a regular pentagon together with a sixth line parallel to one of the five lines forming at least six congruent triangles.

We conclude that in an arrangement on six lines which is f -optimal, every sub-arrangement on five lines contains at most four good triangles. With this in mind, by applying Lemma 2.3 (with $k = 5$), we obtain the desired $f(6) = 8$. Lemmas 2.1 and 2.2 yield the bounds on $g(6)$.

Theorem 3.5 (i) and the Star of David (which proves that $6 \in G(6)$) imply that $[0..6] \subset G(6)$. Together with the arrangements from Figure 5, we are done. \square

Among arrangements on six lines bounding exactly six congruent facial triangles, we found three combinatorially non-equivalent ones. (It is unknown whether these are all.)

In the general case, the solution from Figure 5 (b) seems to be unique; see Conjecture 4.1 (which states that $g(6) = 6$) in the final section.

3.4 $\ell = 7$

Theorem 3.7. We have $f(7) = 14$, $9 \leq g(7) \leq 11$, $[0..10] \cup \{14\} \subset F(7)$, and $[0..9] \subset G(7)$.

Proof. By Theorem 3.6, $f(6) = 8$. Thus, by Lemma 2.2 and Lemma 2.4 (with $k = 6$), $f(7) = 14$. For $g(7)$, the lower bound is given by the construction in Figure 7 (a) (by deleting the line marked h), the upper bound by Lemma 2.1.

Since the Star of David is 1-extendable, we have $7 \in G(7)$. Removing the line marked h in Figures 7 (a) and (b) shows that $9 \in G(7)$ and $8 \in G(7)$, resp. Thus, $[0..9] \subset G(7)$. By considering seven of the eight lines bounding a regular octagon, we obtain $10 \in F(7)$. Together with Lemma 2.2, we have $[0..10] \cup \{14\} \subset F(7)$. \square

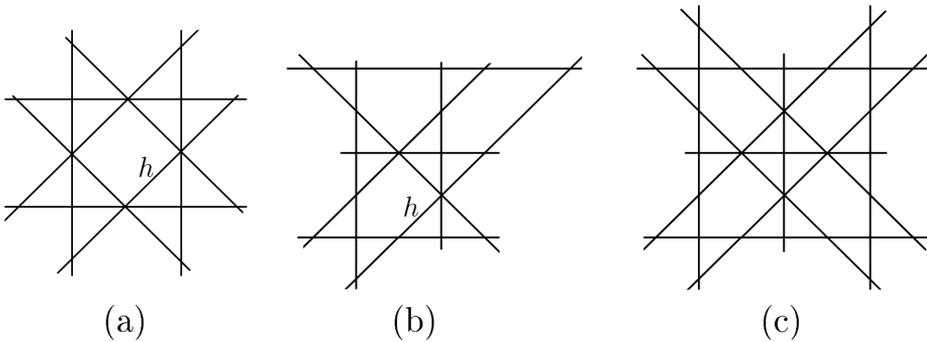


Figure 7: (a) An arrangement proving $12 \in G(8)$. Deleting the line marked h shows that $9 \in G(7)$. (b) An arrangement showing $11 \in G(8)$. Deleting h yields $8 \in G(7)$. (c) This arrangement proves that $g(10) \geq 20$.

3.5 $\ell = 8$

Theorem 3.8. We have $16 \leq f(8) \leq 22$, $12 \leq g(8) \leq 15$, $[0..16] \setminus \{13\} \subset F(8)$ and $[0..12] \subset G(8)$.

Proof. Lemma 2.2 implies the lower bound for $f(8)$, Theorem 3.7 and Lemma 2.4 (with $k = 7$) the upper bound. For $g(8)$, the lower bound is given by the arrangement from Figure 7 (a), the upper bound by Lemma 2.1.

Figures 7 (a) and (b) show that $\{11, 12\} \subset G(8)$. This, Theorem 3.7 and the fact that the arrangement from Figure 7 (a) minus the line marked h is 1-extendable (which proves that $10 \in G(8)$) yield $[0..12] \subset G(8)$. Applying Lemmas 2.2 and 2.4, we obtain $[0..16] \setminus \{13\} \subset F(8)$. \square

3.6 $9 \leq \ell \leq 12$

As the techniques for proving the following results are very similar to what has been shown, we skip them. A notable exception is the construction from Figure 8.

Theorem 3.9. *We have*

$$18 \leq f(9) \leq 33, 15 \leq g(9) \leq 21, [0..18] \subset F(9), [0..15] \subset G(9),$$

$$21 \leq f(10) \leq 48, 20 \leq g(10) \leq 26, [0..21] \subset F(10), [0..20] \subset G(10),$$

$$26 \leq f(11) \leq 66, 23 \leq g(11) \leq 33, [0..26] \subset F(11), [0..23] \subset G(11),$$

and

$$32 \leq f(12) \leq 88, 26 \leq g(12) \leq 39, [0..28] \cup \{32\} \subset F(12), [0..26] \subset G(12).$$

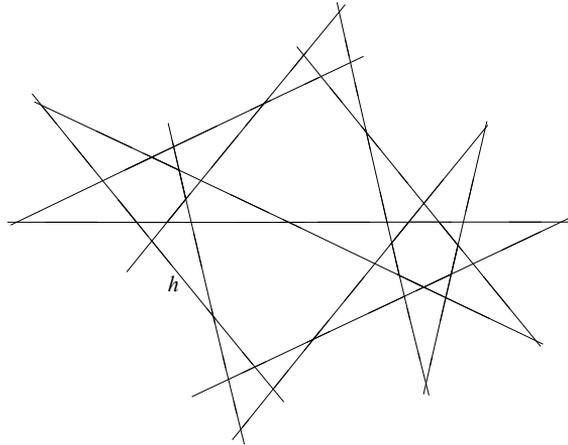


Figure 8: An arrangement proving $f(11) \geq 26$. It is obtained by joining the two arrangements from Figures 1 and 5 (b) with the technique described in the proof of Proposition 3.1, i.e. such that the two arrangements share a pair of lines (forming the same good angle) in the new arrangement. Deleting the line marked h , one obtains $f(10) \geq 21$. By completing the left regular heptagon, we obtain $f(12) \geq 32$.

3.7 Summary

Consider $\ell \leq 12$ lines in the Euclidean plane, and let $f(\ell)$ and $g(\ell)$ be defined as in the Introduction. Then we have the following bounds.

Table 2: Bounding $f(\ell)$ and $g(\ell)$ for $\ell \leq 12$.

ℓ	3	4	5	6	7	8	9	10	11	12
$f(\ell) \geq$	1	2	5	8	14	16	18	21	26	32
$f(\ell) \leq$	1	2	5	8	14	22	33	48	66	88
$g(\ell) \geq$	1	2	5	6	9	12	15	20	23	26
$g(\ell) \leq$	1	2	5	7	11	15	21	26	33	39

We were also able to prove that $f(13) \geq 37, f(14) \geq 44, f(15) \geq 50, f(16) \geq 56,$ and $f(17) \geq 61$.

4 Conjectures

Conjecture 4.1. $g(6) = 6$.

If Conjecture 4.1 is true, we would have $g(7) \leq 10$.

Conjecture 4.2. $g(7) = 9$.

If Conjecture 4.2 is true, we would have $g(8) \leq 14$.

Conjecture 4.3. *The f -optimal arrangements on 6 and 7 lines (consider Figures 5 (b) and 1, resp.) are g -unique.*

Conjecture 4.4. *(a) $F(7)$ is not complete, but (b) for every ℓ , $G(\ell)$ is complete.*

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