

The distinguishing index of the Cartesian product of countable graphs

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Abstract

The *distinguishing index* $D'(G)$ of a graph G is the least cardinal d such that G has an edge colouring with d colours that is preserved only by the trivial automorphism.

We derive some bounds for this parameter for infinite graphs. In particular, we investigate the distinguishing index of the Cartesian product of countable graphs.

Finally, we prove that $D'(K_2^{\aleph_0}) = 2$, where $K_2^{\aleph_0}$ is the infinite dimensional hypercube.

Keywords: Distinguishing index, automorphism, infinite graph, edge colouring, infinite dimensional hypercube.

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1 Introduction

Albertson and Collins [1] introduced the (*vertex-*)*distinguishing number* $D(G)$ of a graph G as the least cardinal d such that G has a labelling with d labels that is only preserved by the trivial automorphism. This concept has spawned numerous papers, mostly on finite graphs. But countable infinite graphs have also been investigated with respect to the distinguishing number; see [12], [13], and [14]. For graphs of higher cardinality, see [8]. The corresponding notion for endomorphisms instead of automorphisms is investigated in [5].

Let us consider now any *edge colouring* of a graph G ; it is merely a function $f : E(G) \rightarrow C$ which labels each edge of G with a *colour* from some set C . Given a graph

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G with an edge colouring f , we say that a graph automorphism $\varphi : V(G) \rightarrow V(G)$ of G preserves the edge colouring f if $f(xy) = f(\varphi(x)\varphi(y))$ for every edge $xy \in E(G)$; if, on the other hand, there is an edge xy such that $f(xy) \neq f(\varphi(x)\varphi(y))$, then we say that φ is broken by xy . It is easy to see that there is, for every connected graph $G \neq K_2$, an edge colouring of G which is preserved only by the trivial automorphism of G , i.e., only by the identity $\text{id}_G : V(G) \rightarrow V(G)$: Merely choose different colours for different edges. The distinguishing index $D'(G)$ of a graph G is the least cardinal d such that G has an edge colouring with d colours that is only preserved by the trivial automorphism. Obviously for K_2 the distinguishing index is not defined and it is the only such connected graph.

For finite graphs this concept is investigated by Kalinowski and Piłśniak in [9] and by Piłśniak in [11]. In [2], the following general upper bound was proved.

Theorem 1.1. *Let G be a connected, infinite graph such that the degree of every vertex of G is not greater than Δ . Then $D'(G) \leq \Delta$.*

A graph G is said to be prime with respect to the Cartesian product if whenever $G \cong G_1 \square G_2$, then either G_1 or G_2 is the graph consisting of a single vertex. It is well known (see [6]) that if G is connected, then G has a unique prime factorization, i.e.,

$$G \cong G_1 \square G_2 \square \cdots \square G_t$$

such that for $1 \leq i \leq t$, G_i is prime. Two graphs G and H are called relatively prime if K_1 is the only common factor of G and H . About forty-five years ago Imrich and Miller independently proved the following theorem – see Thm. 6.10 in [6].

Theorem 1.2. *If G is connected and $G = G_1 \square G_2 \square \cdots \square G_r$ is its prime decomposition, then every automorphism of G is generated by the automorphisms of the factors and the transpositions of isomorphic factors.*

A basic fact, which is a reformulation of the above theorem for $r = 2$ and which is used frequently in this paper, is:

If φ is an automorphism of the Cartesian product $G_1 \square G_2$ of two connected relatively prime graphs, then there are automorphisms φ_i of G_i , $i = 1, 2$, such that $\varphi(v_1, v_2) = (\varphi_1(v_1), \varphi_2(v_2))$ for all $(v_1, v_2) \in V(G_1 \square G_2)$.

In this case we write $\varphi = (\varphi_1, \varphi_2)$ for short and we note that φ is non-trivial if and only if at least one of φ_1 and φ_2 is non-trivial.

An asymmetric graph has only one automorphism, the trivial automorphism. We now state an easy corollary of these properties and definitions for product graphs with distinguishing index 1.

Proposition 1.3. *Let G be the Cartesian product of two graphs G_1 and G_2 . Then*

$$D'(G_1 \square G_2) = 1$$

if and only if G_1 and G_2 are relatively prime and both are asymmetric graphs.

The aim of this paper is to present new results for the distinguishing index of the Cartesian product of infinite graphs. Most graphs in this document are countable, i.e., finite or denumerable; numbers used are either finite or \aleph_0 .

Subgraphs of the Cartesian product $G_1 \square G_2$ of the form $G_1 \square \{v\}$ (for any $v \in V(G_2)$) are isomorphic to G_1 and are called G_1 -layers of $G_1 \square G_2$. The G_2 -layers of $G_1 \square G_2$ are defined similarly.

The distinguishing index of the Cartesian product of finite graphs is investigated in [4] where the authors prove, amongst others, a result which will be useful in the next section and which we now record as

Theorem 1.4. *Let G be a connected finite graph and $k \geq 2$. Then $D'(G^k) = 2$ with the only exception: $D'(K_2^2) = 3$.*

2 The distinguishing index of the Cartesian product

First we consider the Cartesian product of two denumerable graphs with infinite edge sets.

Lemma 2.1. *Let G_1 and G_2 be two connected relatively prime denumerable graphs. Then $D'(G_1 \square G_2) \leq 2$.*

Proof. We start by labelling the edges of G_1 with e_1, e_2, \dots and those of G_2 with f_1, f_2, \dots . This is possible since both edge sets have to be denumerable. Note that these labellings effectively *order* the edges of these graphs. We can now easily describe the required edge distinguishing colouring in colours 1 and 2:

Colour the first (in terms of the above ordering) k edges of the k 'th layer of G_1 and the first k edges of the k 'th layer of G_2 with 1; colour all other edges with 2. Recall that every edge in $G_1 \square G_2$ lies in a G_1 -layer or a G_2 -layer; hence this process colours indeed all edges of $G_1 \square G_2$. Using the labels, this means that the edges corresponding to the edges $\{e_1, e_2, \dots, e_k\}$ of G_1 in the k 'th G_1 -layer and the edges corresponding to the edges $\{f_1, f_2, \dots, f_k\}$ of G_2 in the k 'th G_2 -layer, for all $k = 1, 2, \dots$, are coloured 1 and all other edges are coloured 2.

Now consider, if possible, any non-trivial automorphism $\varphi = (\varphi_1, \varphi_2)$ of $G_1 \square G_2$ which preserves the above edge colouring of $G_1 \square G_2$. Since every two different G_1 -layers have different numbers of edges coloured with 1, the automorphism φ_2 of G_2 must be trivial. Similarly, φ_1 must be trivial. Hence φ is the trivial automorphism, proving that for every non-trivial automorphism φ of $G_1 \square G_2$ there is an edge e of $G_1 \square G_2$ for which e and $\varphi(e)$ are coloured differently. \square

The same result was obtained for the distinguishing number of two connected relatively prime denumerable graphs by Imrich and Klavžar in [7]. Recently it was shown by Estaji, Imrich, Kalinowski, Pilśniak and Tucker in [3] that the condition that the two graphs are relatively prime can be omitted.

Note that Lemma 2.1 assures us that $D'(G_1 \square G_2)$ is at most two irrespective of the values of $D'(G_1)$ and $D'(G_2)$. Next we consider the case in which both G_1 and G_2 of orders being any cardinals and with finite values for the distinguishing index.

Lemma 2.2. *Suppose G_1 and G_2 are connected relatively prime graphs with finite distinguishing indexes. If $D'(G_i) \leq k_i$, $i = 1, 2$, then $D'(G_1 \square G_2) \leq \max\{k_1, k_2\}$.*

Proof. Since $D'(G_i) \leq k_i$, $i = 1, 2$, there are, with $k = \max\{k_1, k_2\}$, edge colourings f_1 of G_1 and f_2 of G_2 using the colours $1, 2, \dots, k$ which are distinguishing colourings of G_1 and G_2 respectively. In order to prove now that $D'(G_1 \square G_2) \leq k$, we again use the notion of a “first” layer through a labelling of the vertices (which here is not explicitly chosen or named). Hence consider the function $f : E(G_1 \square G_2) \rightarrow \{1, 2, \dots, k\}$ defined by
1) $f((v_1, w)(v_2, w)) = f_1(v_1 v_2)$ for edges of the first G_1 -layer and

- 2) $f((v, w_1)(v, w_2)) = f_2(w_1 w_2)$ for edges of the first G_2 -layer and
 3) $f(e) = 1$ for all remaining edges.

Consider any non-trivial automorphism $\alpha = (\alpha_1, \alpha_2)$ of $G_1 \square G_2$ with α_1 a non-trivial automorphism of G_1 or α_2 a non-trivial automorphism of G_2 . Assume that the first is true (for G_1): Then, since f_1 is a distinguishing colouring of the first G_1 -layer, there is an edge e of G_1 such that $f_1(e) \neq f_1(\alpha_1(e))$. Now, if α_2 does not move the first layer, then this edge (considered as an edge of $G_1 \square G_2$) is an edge of the required kind in the first G_1 -layer. On the other hand, if α_2 does move the first layer to another layer, we can remark, since $f_1(e) \neq f_1(\alpha_1(e))$, that at least one of $f_1(e)$ and $f_1(\alpha_1(e))$ is different from 1 so that this edge is moved by α_2 to an edge in another layer which has colour 1 by 3) above.

A similarly argument holds if the second is true (for G_2) – merely interchange the roles of G_1 and G_2 (and their colourings and automorphisms) in the above argument.

Hence we are assured that all non-trivial automorphisms of $G_1 \square G_2$ are broken by the colouring f . \square

Observe, that $D'(G_1 \square G_2)$ can be arbitrary large, for instance if G_1 is isomorphic to P_3 and G_2 is isomorphic to an infinite ray with many (but finitely many) leaves adjacent to its first vertex.

In our next result we prove that if G_1 satisfies $D'(G_1) = \aleph_0$ and the graph G_2 is finite (so that, in particular $D'(G_2)$ is finite), then $D'(G_1 \square G_2) = \aleph_0$.

Lemma 2.3. *Suppose G_1 and G_2 are connected relatively prime graphs with $D'(G_1) = \aleph_0$ and G_2 is finite. Then $D'(G_1 \square G_2) = \aleph_0$.*

Proof. Suppose, for a proof by contradiction, that $D'(G_1 \square G_2)$ is finite. Since G_2 is a finite graph, there are finite values for $\|G_2\|$, the number of edges of G_2 , and $D'(G_2)$ too. Hence we can choose a positive integer k such that each of these three numbers is at most k .

Since $D'(G_1 \square G_2) \leq k$, there is a k -distinguishing edge colouring f of the edges of $G_1 \square G_2$. Furthermore, since $D'(G_1) = \aleph_0$, there exists, for every positive integer t , a non-trivial automorphism α_t of G_1 which needs at least $t + 1$ colours to break it. So if $t \geq k$, the colouring by f of any layer of G_1 induces a colouring on G_1 which cannot be broken by the automorphism α_t of G_1 . Since there are infinitely many such automorphisms, we may assume without loss of generality that $\alpha_s \neq \alpha_t$ when $s \neq t$.

Now consider non-trivial automorphisms of $G_1 \square G_2$ of the form $\alpha = (\alpha_t, \text{id}_{G_2})$ (for some $t \geq k$). For each such t , and each edge vw of G_1 (which we can consider as an edge of any G_1 -layer of $G_1 \square G_2$), we have that $f(vw) = f(\alpha_t(v)\alpha_t(w))$, i.e., these automorphisms of $G_1 \square G_2$ are not broken by edges in layers of G_1 .

The automorphisms α of the above form should therefore be broken by edges of layers of G_2 . But this means that, for each $t \geq k$, for at least one edge xy of the G_2 -layer determined by a vertex $v \in V(G_1)$, we have that $f(xy)$ in this layer is different from $f(\alpha_t(x)\alpha_t(y))$ in the G_2 -layer determined by $\alpha_t(v) \in V(G_1)$. Since there are infinitely many G_2 -layers, this requires infinitely many different colourings of G_2 . However, there are at most $k^{\|G_2\|}$ different colourings of G_2 -layers. Hence the colouring f cannot break all the infinitely many automorphisms described above. \square

As a consequence of the above three lemmas we immediately obtain the following characterisation.

Theorem 2.4. *If G_1 and G_2 are connected relatively prime countable graphs, then $D'(G_1 \square G_2)$ is infinite if and only if for some $i \in \{1, 2\}$ we have that $D'(G_i)$ is infinite while for $j \neq i$ we have that G_j is finite. \square*

Now we consider a graph which is the Cartesian power G^k of a denumerable graph G . For a finite graph G , the distinguishing number of the Cartesian power of G is considered in [4]. Here we prove a result for graphs G with a finite number of prime factors (counted with their multiplicities). We begin with a result for prime graphs.

Lemma 2.5. *Let $k \geq 2$ be an integer. If a connected denumerable graph G is prime with respect to the Cartesian product, then $D'(G^k) = 2$.*

Proof. If $k = 2$, the proof is similar to the proof of Lemma 2.1. Indeed, denote $G^2 = G_1 \square G_2$, where G_1, G_2 are isomorphic to G . Using an analogous proof technique but colouring distinct even numbers of edges of each G_1 -layer with red and distinct odd numbers of edges of each G_2 -layer with red will also take care of the additional automorphisms generated by the isomorphism between G_1 and G_2 .

Now we show that $D'(G \square H) = 2$ if $D'(H) = 2$ and G is prime. In particular, if we consider $H = G^{k-1}$ then we obtain the thesis by induction. Namely, let f be a distinguishing colouring of H with two colours. We define a colouring of $G \square H$ as follows: One H -layer is given the colouring f , hence all automorphisms of this H -layer are broken. We colour another H -layer completely blue and all remaining H -layers we colour with distinct numbers of red edges different from the number of red edges in f . Hence all automorphisms of G are broken. If G' isomorphic with G is a factor of H , then we have additional automorphisms, generated by interchanging of G and G' . To break them, we colour each G -layer red. Then every G' -layer contained in a blue H -layer is completely blue, so it cannot be interchanged with G . In this way we break all nontrivial automorphisms of $G \square H$ with two colours if $D'(H) = 2$ and G is prime. \square

The above proof is analogous to the proof of a similar result in [7]. Observe that $D'(G \square H) = 2$ if $D'(H) = 2$ and G is prime, also if G is finite.

Theorem 2.6. *Let $k \geq 2$ be an integer and G be a connected denumerable graph with the prime factor decomposition $G = G_1 \square \dots \square G_r$, where G_1, \dots, G_r are not necessarily distinct. Then $D'(G^k) = 2$.*

Proof. If G is prime, the claim follows from Lemma 2.5. If G is not prime, we consider the prime factorization $G = G_1 \square \dots \square G_r$ and apply Lemma 2.5 to every infinite factor (G has at least one infinite prime factor). Moreover, we can use Theorem 1.4 for every finite factor. The result then follows from Lemma 2.2 unless $G = K_2 \square H$ and $k = 2$, where H is an infinite graph relatively prime with K_2 . But we already know that $D'(H^2) = 2$ due to the above arguments, so let f be a distinguishing colouring of H^2 with two colours. We then define a colouring of G^2 in terms of its four H^2 -layers as follows: One H^2 -layer is given the colouring f , hence all automorphisms of this H^2 -layer are broken. The three remaining H^2 -layers are coloured with distinct numbers of red edges (while all remaining edges are blue), hence all automorphisms of G^2 are broken. \square

We say the G has infinite diameter if there are vertices of arbitrarily large distance. Such a situation occurs in particular in any weak Cartesian product G of infinitely many non-trivial factors (finite or infinite). Hence the above theorem immediately implies the following.

Corollary 2.7. *Let $k \geq 2$ be an integer and let G be a connected denumerable graph with finite diameter. Then $D'(G^k) = 2$.*

3 The distinguishing index of the infinite hypercube

The situation is quite different when we have infinitely many factors in the Cartesian power – consider for example the *infinite dimensional hypercube* $K_2^{\aleph_0}$. This (uncountable) graph has vertices represented by (denumerable) sequences of 0's and 1's and two vertices are adjacent whenever their binary sequences differ in exactly one entry. This graph also has uncountably many connected components, each a countable graph, which are pairwise isomorphic. The automorphism group of $K_2^{\aleph_0}$ is well described (see [10]). Using this information, we are now ready to prove

Theorem 3.1. *Let $K_2^{\aleph_0}$ be the infinite dimensional hypercube. Then $D'(K_2^{\aleph_0}) = 2$.*

Proof. We first construct an asymmetric spanning tree and then show how it can be used to prove the existence of an asymmetric spanning subgraph in every component of $K_2^{\aleph_0}$; these subgraphs will be constructed in such a way that different components have non-isomorphic subgraphs. Towards the end of the proof, we shall show how they can be exploited to break all non-trivial automorphisms of the hypercube $K_2^{\aleph_0}$.

It is convenient to describe the required asymmetric subgraphs by first handling the connected component C^0 in which all sequences have only finitely many 1's (and therefore an infinite tail of 0's). First we build an asymmetric tree T , which is a spanning subgraph of C^0 , as follows:

Take $(0, 0, 0, 0, \dots)$ and let it be the central vertex. Then add $(1, 0, 0, 0, \dots)$, and the edge between it and the central vertex, to form the first branch of the tree. Next take $(0, 1, 0, 0, 0, \dots)$ and $(1, 1, 0, 0, 0, \dots)$ and the path between them and the central vertex to form the second branch of the tree. The i 'th branch of this tree will therefore be the path on the central vertex and $(0^{i-1}, 1, 0, 0, 0, \dots)$, $(0^{i-2}, 1, 1, 0, 0, 0, \dots)$, \dots and will have length 2^{i-1} . All these binary sequences have 1 on the i 'th entry and if we restricted them to the first $i - 1$ entries, then we obtain the binary-reflected Gray code list with $i - 1$ bits. It can be generated recursively from the list for $i - 2$ bits by reflecting the list (i.e. listing the entries in reverse order), concatenating the original list with the reversed list, prefixing the entries in the original list with 0, and then prefixing the entries in the reflected list with 1. In particular, the last vertex of the i 'th branch has the code $(1, 0^{i-2}, 1, 0, 0, 0, \dots)$, and the last but one has the code $(1, 0^{i-3}, 1, 1, 0, 0, 0, \dots)$.

Note that all branches of T are of different length, which ensures us that T is asymmetric, and note that it is a spanning tree of the component C^0 . So it means that we can easily distinguish the weak Cartesian product of \aleph_0 copies of K_2 by two colors: Namely we colour all the edges of T with one colour and the remaining edges with the second colour.

Now we would like to distinguish the Cartesian product of \aleph_0 copies of K_2 by two colours. Consider any sequence $\mathbf{x} = (x_1, x_2, \dots)$ of 0's and 1's and suppose it is in the connected component C of the hypercube $K_2^{\aleph_0}$. Since C is isomorphic to C^0 , we can find a copy of T , say T^C , in C . Now we use \mathbf{x} and T^C to create a spanning subgraph $T_{\mathbf{x}}^C$ of C by adding edges to T^C as follows:

For every positive integer i we add the edge of $K_2^{\aleph_0}$ between the endvertex of the i 'th branch and the last but one vertex of the $(i + 1)$ 'th branch of T^C to this tree if and only if

$x_i = 1$. We remark that this edge is indeed in $K_2^{\aleph_0}$ since the binary sequences representing these vertices in C^0 differ in exactly one entry, namely the $(i+1)$ 'th entry, and therefore the same is true in the isomorphic copy T^C of T . Note also that the choice of the added edges ensures us that $T_{\mathbf{x}}$ is not isomorphic to $T_{\mathbf{x}'}$ whenever $\mathbf{x} \neq \mathbf{x}'$. Since there are uncountably many sequences \mathbf{x} , we thus have uncountably many pairwise non-isomorphic subgraphs all of which are asymmetric.

Finally we prove, using these subgraphs of the components of $K_2^{\aleph_0}$, that the infinite hypercube is 2-distinguishable. Consider the following colouring f of the edges of $K_2^{\aleph_0}$: Colour, for each component C of $K_2^{\aleph_0}$ and some fixed choice of a vertex \mathbf{x} of C , all the edges of the spanning subgraph $T_{\mathbf{x}}^C$ with 1; colour all the other edges of $K_2^{\aleph_0}$ with 2. Then consider any automorphism α of $K_2^{\aleph_0}$. Since isomorphisms, and thus α , preserve connectivity, α has to take every component C of $K_2^{\aleph_0}$ to a component C' of $K_2^{\aleph_0}$. But, if $C \neq C'$, then the asymmetric spanning subgraphs $T_{\mathbf{x}}^C$ and $T_{\mathbf{x}'}^{C'}$ of C and C' are not isomorphic (because $\mathbf{x} \neq \mathbf{x}'$), hence the colouring f breaks α . \square

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