

Reconfiguring vertex colourings of 2-trees

Michael Cavers , Karen Seyffarth

*Department of Mathematics and Statistics, University of Calgary,
Calgary, AB T2N 1N4 Canada*

Received 4 October 2018, accepted 8 October 2019, published online 16 December 2019

Abstract

Let H be a graph and let $k \geq \chi(H)$ be an integer. The k -colouring graph of H , denoted $G_k(H)$, is the graph whose vertex set consists of all proper k -vertex-colourings (or simply k -colourings) of H using colours $\{1, 2, \dots, k\}$; two vertices of $G_k(H)$ are adjacent if and only if the corresponding k -colourings differ in colour on exactly one vertex of H . If $G_k(H)$ has a Hamilton cycle, then H is said to have a *Gray code* of k -colourings, and the *Gray code number* of H is the least integer $k_0(H)$ such that $G_k(H)$ has a Gray code of k -colourings for all $k \geq k_0(H)$. Choo and MacGillivray determine the Gray code numbers of trees. We extend this result to 2-trees. A 2-tree is constructed recursively by starting with a complete graph on three vertices and connecting each new vertex to an existing clique on two vertices. We prove that if H is a 2-tree, then $k_0(H) = 4$ unless H is isomorphic to the join of a tree T and a vertex u , where T is a star on at least three vertices, or the bipartition of T has two even parts; in these cases, $k_0(H) = 5$.

Keywords: 2-trees, graph colouring, Gray codes, Hamilton cycles, reconfiguration problems.

Math. Subj. Class.: 05C15, 05C45

Rekonfiguracija vozliščnih barvanj 2-dreves

Michael Cavers, Karen Seyffarth

*Department of Mathematics and Statistics, University of Calgary,
Calgary, AB T2N 1N4 Canada*

Prejeto 4. oktobra 2018, sprejeto 8. oktobra 2019, objavljeno na spletu 16. decembra 2019

Povzetek

Naj bo H graf in naj bo $k \geq \chi(H)$ celo število. Graf k -barvanj grafa H , označen z $G_k(H)$, je graf, katerega množica vozlišč sestoji iz vseh pravih k -vozliščnih barvanj (ali preprosto k -barvanj) grafa H z barvami $\{1, 2, \dots, k\}$; vozlišči grafa $G_k(H)$ sta sosednji, če in samo če se ustrezni k -barvanji razlikujeta v barvi natanko enega vozlišča grafa H . Če ima $G_k(H)$ hamiltonski cikel, potem pravimo, da ima H Grayevo kodo iz k -barvanj. Najmanjše celo število $k_0(H)$, pri katerem ima $G_k(H)$ Grayevo kodo iz k -barvanj za vse $k \geq k_0(H)$, imenujemo prag Grayeve kode grafa H . Choo in MacGillivray sta določila pragove Grayevih kod za drevesa. V pričujočem članku razširimo ta rezultat na 2-drevesa. Konstrukcija 2-drevesa poteka rekurzivno: začnemo s polnim grafom na treh vozliščih, potem pa vsako novo vozlišče dodamo neki že obstoječi klikli na dveh vozliščih. Dokažemo, da če je H 2-drevo, potem je $k_0(H) = 4$, razen če je H izomorfen spoju drevesa T in vozlišča u , kjer je T zvezda na najmanj treh vozliščih, ali če se T deli na dva enako velika dela; v teh dveh primerih je $k_0(H) = 5$.

Ključne besede: 2-drevesa, barvanje grafa, Grayeve kode, hamiltonski cikli, rekonfiguracijski problemi.

Math. Subj. Class.: 05C15, 05C45
