

On the parameters of intertwining codes*

Stephen P. Glasby, Cheryl E. Praeger

*Centre for Mathematics of Symmetry and Computation,
University of Western Australia, 35 Stirling Highway, Crawley 6009, Australia*

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Abstract

Let F be a field and let $F^{r \times s}$ denote the space of $r \times s$ matrices over F . Given equinumerous subsets $\mathcal{A} = \{A_i \mid i \in I\} \subseteq F^{r \times r}$ and $\mathcal{B} = \{B_i \mid i \in I\} \subseteq F^{s \times s}$ we call the subspace $C(\mathcal{A}, \mathcal{B}) := \{X \in F^{r \times s} \mid A_i X = X B_i \text{ for } i \in I\}$ an *intertwining code*. We show that if $C(\mathcal{A}, \mathcal{B}) \neq \{0\}$, then for each $i \in I$, the characteristic polynomials of A_i and B_i share a nontrivial factor. We give an exact formula for $k = \dim(C(\mathcal{A}, \mathcal{B}))$ and give upper and lower bounds. This generalizes previous work. Finally we construct intertwining codes with large minimum distance when the field is not ‘too small’. We give examples of codes where $d = rs/k = 1/R$ is large where the minimum distance, dimension, and rate of the linear code $C(\mathcal{A}, \mathcal{B})$ are denoted by d , k , and $R = k/rs$, respectively.

Keywords: Linear code, dimension, distance.

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E-mail addresses: Stephen.Glasby@uwa.edu.au (Stephen P. Glasby), Cheryl.Praeger@uwa.edu.au (Cheryl E. Praeger)

O parametrih prepletnih kod*

Stephen P. Glasby, Cheryl E. Praeger

*Centre for Mathematics of Symmetry and Computation,
University of Western Australia, 35 Stirling Highway, Crawley 6009, Australia*

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Povzetek

Naj bo F obseg, $F^{r \times s}$ pa prostor $r \times s$ matrik nad F . Če sta podani enako močni podmnožici $\mathcal{A} = \{A_i \mid i \in I\} \subseteq F^{r \times r}$ in $\mathcal{B} = \{B_i \mid i \in I\} \subseteq F^{s \times s}$, potem pravimo podprostoru $C(\mathcal{A}, \mathcal{B}) := \{X \in F^{r \times s} \mid A_i X = X B_i \text{ za } i \in I\}$ *prepletna koda*. Pokažemo, da če je $C(\mathcal{A}, \mathcal{B}) \neq \{0\}$, potem si, za vsak $i \in I$, karakteristična polinoma matrik A_i in B_i delita netrivialen faktor. Izpeljemo natančno formulo za $k = \dim(C(\mathcal{A}, \mathcal{B}))$ in podamo zgornjo in spodnjo mejo. To je posplošitev prejšnjih rezultatov. Konstruiramo prepletne kode z veliko minimalno razdaljo v primerih, ko obseg ni ‘premajhen’. Podamo primere prepletov, kjer je $d = rs/k = 1/R$ velik, pri čemer so d , k , in $R = k/rs$ oznake za minimalno razdaljo, dimenzijo in velikost prepleta, v tem vrstnem redu.

Ključne besede: Linearna koda, dimenzija, razdalja.

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E-poštna naslova: Stephen.Glasby@uwa.edu.au (Stephen P. Glasby), Cheryl.Praeger@uwa.edu.au (Cheryl E. Praeger)