

A characterization of graphs with disjoint total dominating sets*

Michael A. Henning[†]

*Department of Pure and Applied Mathematics, University of Johannesburg,
Auckland Park, 2006 South Africa*

Iztok Peterin[‡]

*Faculty of Electrical Engineering and Computer Science, University of Maribor,
Koroška 46, 2000 Maribor, Slovenia*

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Abstract

A set S of vertices in a graph G is a total dominating set of G if every vertex is adjacent to a vertex in S . A fundamental problem in total domination theory in graphs is to determine which graphs have two disjoint total dominating sets. In this paper, we solve this problem by providing a constructive characterization of the graphs that have two disjoint total dominating sets. Our characterization gives an entirely new description of graphs with two disjoint total dominating sets and places them in another context, developing them from four base graphs and applies a sequence of operations from seventeen operations that are independent and necessary to produce all such graphs. We show that every graph with two disjoint total dominating sets can be constructed using this method.

Keywords: Total domination number, disjoint total dominating sets.

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[†]Research supported in part by the South African National Research Foundation and the University of Johannesburg.

[‡]This author is also affiliated with Institute of Mathematics, Physics and Mechanics, Jadranska 19, 1000 Ljubljana, Slovenia. Research supported in part by the Slovenian Research Agency by the projects No. P1-0297 and No. J1-9109.

E-mail addresses: mahenning@uj.ac.za (Michael A. Henning), iztok.peterin@um.si (Iztok Peterin)

1 Introduction

A *dominating set* of a graph G is a set S of vertices of G such that every vertex not in S has a neighbor in S , where two vertices are neighbors if they are adjacent. A *total dominating set* of a graph G with no isolated vertex is a set S of vertices such that every vertex in G has a neighbor in S . Domination and its variations in graphs are now well studied. The literature on this subject has been surveyed and detailed in the two books by Haynes, Hedetniemi, and Slater [6, 7]. For a recent book on total domination in graphs we refer the reader to [13]. A survey of total domination in graphs can also be found in [9].

A classical result in domination theory, due to Ore [14] in 1962, is that every graph with no isolated vertex has two disjoint dominating sets. However, it is not the case that every graph with no isolated vertex can be partitioned into a dominating set and a total dominating set. Henning and Southey [11] showed that every connected graph with minimum degree at least two that is not a cycle on five vertices has a disjoint dominating set and a total dominating set. Further, in [12] they present a constructive characterization of connected graphs of order at least 4 that have a disjoint dominating set and a total dominating set. Disjoint dominating and total dominating sets in graphs are studied further, for example, in [10]. A characterization of graphs with disjoint dominating and paired-dominating sets is characterized in [15].

It remains, however, an outstanding problem to determine which graphs have two disjoint total dominating sets. Zelinka [16] in 1989 showed that no minimum degree condition in a graph is sufficient to guarantee that there exist two disjoint total dominating sets in the graph. Heggernes and Telle [8] showed that the decision problem to decide for a given graph G if it has two disjoint total dominating sets is NP-complete, even for bipartite graphs. Sufficient conditions for a graph to have two disjoint total dominating sets were obtained by Delgado, Desormeaux, and Haynes [4], but the authors in [4] were not able to characterize such graphs. Cubic graphs that have two disjoint total dominating sets were recently studied by Desormeaux, Henning and Haynes [5]. In particular, they show that cubic graphs that are 5-chordal or claw-free (we do not define these concepts here) can be partitioned into two total dominating sets.

The *total domatic number* $tdom(G)$ of G is the maximum number of disjoint total dominating sets [3]. This can also be considered as a coloring of the vertices such that every vertex has a neighbor of every color (and has been called the *coupon coloring problem* [2]). Recent work on the total domatic number can be found, for example, in [1, 5]. The fundamental problem in total domination theory in graphs of determining which graphs have two disjoint total dominating sets can be phrased as follows: Determine which graphs G satisfy $tdom(G) \geq 2$. We call a graph a *TDP-graph* (standing for “total dominating partitionable graph”) if its vertex set can be partitioned into two total dominating sets; that is, a graph G is a TDP-graph if and only if $tdom(G) \geq 2$.

In this paper, we provide a constructive characterization of the graphs that have two disjoint total dominating sets, or, equivalently, a characterization of the TDP-graphs. We describe a procedure to build TDP-graphs in terms of a 2-coloring of the vertices that indicate the role each vertex plays in the sets associated with the two disjoint total dominating sets. We show that the resulting family we construct, starting from four initial base graphs and applying one of seventeen operations to extend graphs in the family to larger graphs, is precisely the class of all TDP-graphs.

Our characterization provides a method for creating the TDP-graphs using a finite set of precise operations. The construction places the TDP-graphs in another context, devel-

oping them from four base graphs and applying a sequence of operations from seventeen operations that are independent and necessary to produce a TDP-graph; that is, we show that this method produces precisely the family of TDP-graphs in that every graph generated by this method/algorithm is a TDP-graph and further every TDP-graph can be constructed in this way.

We remark that this procedure does not solve the decision problem to decide if a given graph has two disjoint total dominating sets in polynomial time. If one follows the steps in the proof of Section 4, one does indeed obtain an algorithm for this decision problem. However, this algorithm is far from polynomial time complexity. In particular, the first step of this algorithm is to discard some edges in order to obtain so-called sparse TDP-graph. Unfortunately, the proof does not provide those edges and this already spoils the time complexity.

1.1 Notation

For notation and graph theory terminology we generally follow [13]. All graphs in this paper are finite and simple, without loops or multiple edges. The *order* of a graph G is denoted by $n(G) = |V(G)|$, and the *size* of G by $m(G) = |E(G)|$. We denote the *degree* of a vertex v in the graph G by $d_G(v)$. The maximum (minimum) degree among the vertices of G is denoted by $\Delta(G)$ ($\delta(G)$, respectively). The *open neighborhood* of v is $N_G(v) = \{u \in V(G) \mid uv \in E(G)\}$. For a set $S \subseteq V(G)$, its *open neighborhood* is the set $N_G(S) = \bigcup_{v \in S} N_G(v)$. For subsets X and Y of vertices of G , we denote the set of edges with one end in X and the other end in Y by $[X, Y]$. For a set $S \subseteq V(G)$, the subgraph induced by S is denoted by $G[S]$. Further, the subgraph of G obtained from G by deleting all vertices in S and all edges incident with vertices in S is denoted by $G - S$; that is, $G - S = G[V(G) \setminus S]$. If $S = \{v\}$, we simply denote $G - \{v\}$ by $G - v$.

The *distance* between two vertices u and v in G , denoted $d_G(u, v)$, is the minimum length of a (u, v) -path in G . By W_{uv} we denote the set of all vertices of G which are closer to u than to v ; that is, $W_{uv} = \{w \mid d_G(w, u) < d_G(w, v)\}$. Symmetrically, W_{vu} is defined. A *block* of a graph G is a maximal connected subgraph of G which has no cut-vertex of its own. A block containing exactly one cut-vertex of G is called an *end-block*. It is well known that any two different blocks of a graph have at most one vertex in common, namely a cut-vertex. Furthermore, a connected graph with at least one cut-vertex has at least two end-blocks. Let X denote the set of cut-vertices of a connected graph G and let Y denote the set of its blocks. The *block graph* of G is a bipartite graph with partite sets X and Y in which a vertex $x \in X$ is adjacent to a vertex $y \in Y$ if x is a vertex of the block y . It is well-known that the block graph of any connected graph is a tree.

A *walk* is a finite, alternating sequence of vertices and edges in which each edge of the sequence joins the vertex that precedes it in the sequence to the vertex that follows it in the sequence. A *non-backtracking walk* is a walk with the additional constraint that no two consecutive edges on the walk are repeated.

Let u be a cut-vertex of a graph G . Let H_1 and H_2 be two vertex disjoint subgraphs of $G - u$ that contain all the components of $G - u$, where each of H_1 and H_2 contain at least one component of $G - u$. We call H_1 and H_2 the *associated subgraphs* of $G - u$. For $i \in [2]$, we denote by H_i^u the subgraph of G induced by $V(H_i) \cup \{u\}$. Further, the vertex in H_1^u named u we rename u' , and the vertex in H_2^u named u we rename u'' in order to distinguish between u, u' and u'' . We use the standard notation $[k] = \{1, \dots, k\}$.

2 The graph family \mathcal{G}

In this section, we construct a graph family \mathcal{G} such that every graph in the family has two disjoint total dominating sets. First, we define a 2-coloring of a graph G as a partition $S = (S_R, S_B)$ of $V(G)$. The color of a vertex v , denoted $\text{color}(v)$, is the letter $X \in \{R, B\}$ such that $v \in S_X$, where “R” and “B” here stand for red and blue, respectively. Thus, our 2-coloring of G is a coloring of the vertices of G , one color to each vertex, using the colors red and blue. We denote by \bar{X} the letter $\bar{X} \in \{R, B\} \setminus \{X\}$, and we call \bar{X} the color different from X . Thus, if $X = R$, then $\bar{X} = B$ while if $X = B$, then $\bar{X} = R$. We denote by (G, S) a graph G with a given 2-coloring S . Our aim is to describe a procedure to build TDP-graphs in terms of 2-colorings. For $i \in [4]$, by a 2-colored G_i , we shall mean the graph G_i and its associated 2-coloring shown in Figure 1. Further, we call each 2-colored G_i a 2-colored base graph.

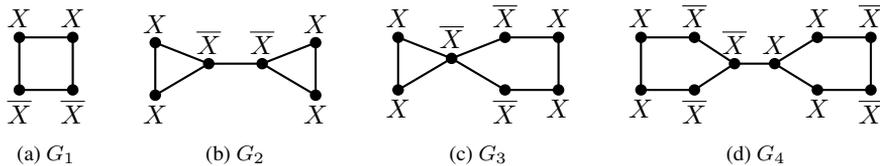


Figure 1: The four 2-colored base graphs G_1, G_2, G_3, G_4 .

Let \mathcal{G} be the minimum family of 2-colored graphs that:

- (i) contains the four 2-colored base graphs; and
- (ii) is closed under the seventeen operations \mathcal{O}_1 through to \mathcal{O}_{17} listed below, which extend a 2-colored graph (G', S') to a new 2-colored graph (G, S) .

In Figures 2–7, the vertices of G' are colored black and the new vertices of G are colored white for illustrative purposes, even though the actual colors of the vertices are indicated by the letters X and \bar{X} .

Operation \mathcal{O}_1 : (G, S) is obtained from (G', S') by adding an edge between two nonadjacent vertices of the same color. See the upper diagram of Figure 2.

Operation \mathcal{O}_2 : (G, S) is obtained from (G', S') by adding an edge between two nonadjacent vertices of different color. See the lower diagram of Figure 2.

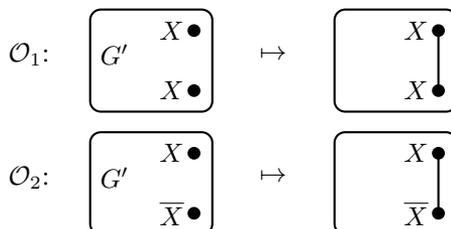


Figure 2: The operations \mathcal{O}_1 and \mathcal{O}_2 .

Operation \mathcal{O}_3 : If u and v are distinct vertices of different color from (G', S') , then (G, S) is obtained from (G', S') by adding a new vertex of any color adjacent to both u and v . See the left diagram in the upper part of Figure 3.

Operation \mathcal{O}_4 : If u and v are distinct vertices of the same color from (G', S') , then (G, S) is obtained from (G', S') by adding adjacent vertices x and y and edges ux and vy with $\text{color}(x) = \text{color}(y) \neq \text{color}(u)$. See the middle diagram in the upper part of Figure 3.

Operation \mathcal{O}_5 : If u and v are distinct vertices of different color from (G', S') , then (G, S) is obtained from (G', S') by adding adjacent vertices x and y and edges ux and vy with $\text{color}(x) = \text{color}(u) \neq \text{color}(y)$. See the right diagram in the upper part of Figure 3.

Operation \mathcal{O}_6 : If u and v are distinct vertices of the same color from (G', S') , then (G, S) is obtained from (G', S') by adding a path xyz with $\text{color}(y) = \text{color}(z) \neq \text{color}(x) = \text{color}(u)$ and adding edges ux and vz . See the left diagram in the lower part of Figure 3.

Operation \mathcal{O}_7 : If u and v are distinct vertices of the same color from (G', S') , then (G, S) is obtained from (G', S') by adding a path $xyzw$ and edges ux and vw with $\text{color}(x) = \text{color}(w) = \text{color}(u) \neq \text{color}(y) = \text{color}(z)$. See the middle diagram in the lower part of Figure 3.

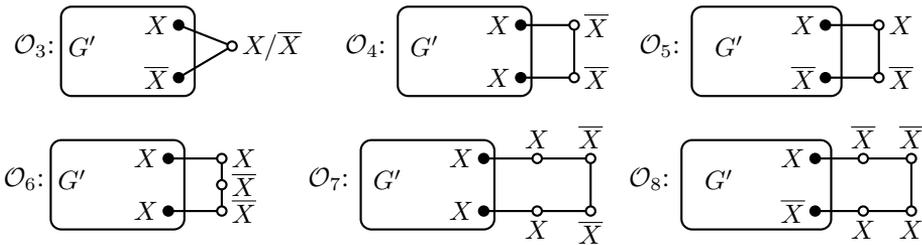


Figure 3: The operations $\mathcal{O}_3 - \mathcal{O}_8$.

Operation \mathcal{O}_8 : If u and v are distinct vertices of different color from (G', S') , then (G, S) is obtained from (G', S') by adding a path $xyzw$ and edges ux and vw with $\text{color}(x) = \text{color}(y) = \text{color}(v) \neq \text{color}(z) = \text{color}(w)$. See the right diagram in the lower part of Figure 3.

Operation \mathcal{O}_9 : If u and v are adjacent vertices of different color from (G', S') , then (G, S) is obtained from (G', S') by subdividing uv with four consecutive vertices x, y, z, w where x is adjacent to u and $\text{color}(u) = \text{color}(z) = \text{color}(w) \neq \text{color}(x) = \text{color}(y)$. See the upper diagram of Figure 4.

Operation \mathcal{O}_{10} : If u and v are adjacent vertices of the same color from (G', S') , then (G, S) is obtained from (G', S') by subdividing uv with four consecutive vertices x, y, z, w where x is adjacent to u and $\text{color}(u) = \text{color}(x) = \text{color}(w) \neq \text{color}(y) = \text{color}(z)$. See the lower diagram of Figure 4.

Operation \mathcal{O}_{11} : If v is a vertex from (G', S') , then (G, S) is obtained from (G', S') by adding an edge xy together with the edges vx and vy where $\text{color}(x) = \text{color}(y) \neq \text{color}(v)$. See the left diagram of Figure 5.

Operation \mathcal{O}_{12} : If v is a vertex from (G', S') , then (G, S) is obtained from (G', S') by

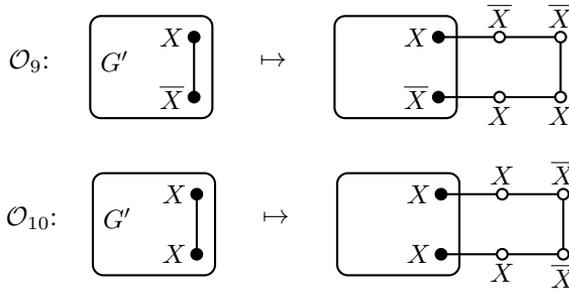


Figure 4: The operations \mathcal{O}_9 and \mathcal{O}_{10} .

adding a path xyz together with the edges vx and vz where $\text{color}(x) = \text{color}(y) \neq \text{color}(z) = \text{color}(v)$. See the middle diagram of Figure 5.

Operation \mathcal{O}_{13} : If v is a vertex from (G', S') , then (G, S) is obtained from (G', S') by adding a path $xyzw$ together with the edges vx and vw where $\text{color}(x) = \text{color}(w) = \text{color}(v) \neq \text{color}(y) = \text{color}(z)$. See the right diagram of Figure 5.

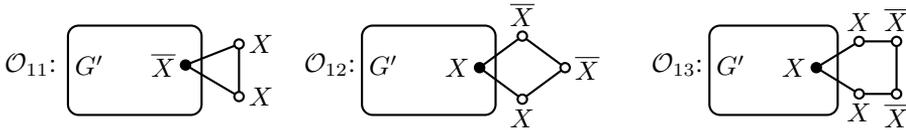


Figure 5: The operations \mathcal{O}_{11} , \mathcal{O}_{12} and \mathcal{O}_{13} .

Operation \mathcal{O}_{14} : If v is a vertex from (G', S') , then (G, S) is obtained from (G', S') by adding a 3-cycle, $xyzx$, together with the edge vx where $\text{color}(x) = \text{color}(v) \neq \text{color}(y) = \text{color}(z)$. See the left diagram of Figure 6.

Operation \mathcal{O}_{15} : If v is a vertex from (G', S') of any color, then (G, S) is obtained from (G', S') by adding a 4-cycle, $xyzwx$, together with the edge vx where $\text{color}(x) = \text{color}(y) \neq \text{color}(z) = \text{color}(w)$. See the middle diagram of Figure 6, where the notation X/\bar{X} means that the vertex can have any color.

Operation \mathcal{O}_{16} : If v is a vertex from (G', S') , then (G, S) is obtained from (G', S') by adding a 5-cycle, $xyzwtx$, together with the edge vx where $\text{color}(x) = \text{color}(y) = \text{color}(t) \neq \text{color}(z) = \text{color}(w) = \text{color}(v)$. See the right diagram of Figure 6.

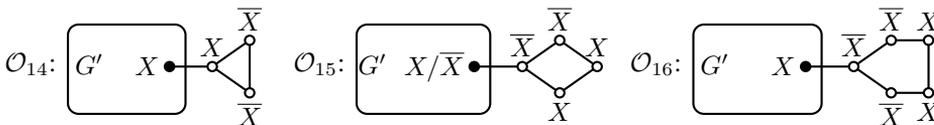


Figure 6: The operations \mathcal{O}_{14} , \mathcal{O}_{15} and \mathcal{O}_{16} .

Operation \mathcal{O}_{17} : If u is a cut-vertex from (G', S') with associated subgraphs H_1^u and H_2^u ,

and in $N_{H_1^u}(u')$ there exists a vertex of the same color as u and in $N_{H_2^u}(u'')$ there exists a vertex of different color as u , then (G, S) is obtained from H_1^u and H_2^u by adding a new vertex v and the edges $u'v$ and vu'' . The color of all vertices from H_1^u remains the same as in G' , $\text{color}(v) = \text{color}(u'') \neq \text{color}(u') = \text{color}(u)$ and the color of all vertices from H_2^u is exchanged with respect to their color in G' . See the diagram of Figure 7, where the notation \bar{A} means that the color of all vertices from the set A in (G', S') is changed in (G, S) .

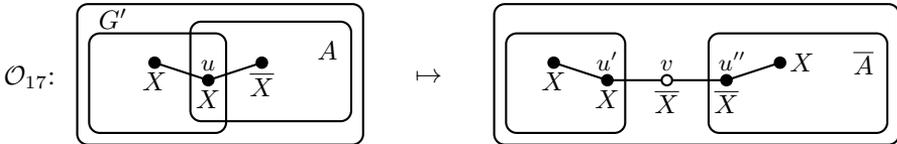


Figure 7: The operation \mathcal{O}_{17} .

We remark that, by definition, all operations \mathcal{O}_3 to \mathcal{O}_{17} produce new vertices. Further, exactly one new vertex created in each of the operations \mathcal{O}_{14} to \mathcal{O}_{16} has degree 3, and all other new vertices created using operations \mathcal{O}_3 to \mathcal{O}_{17} have degree 2 in G . In operations \mathcal{O}_{11} to \mathcal{O}_{13} , if the selected vertex v from (G', S') is a cut-vertex of G' it is also a cut-vertex in G , while if v is not a cut-vertex of G' it becomes a cut-vertex in G . Moreover all operations from \mathcal{O}_{14} to \mathcal{O}_{17} produce new cut vertices. In this sense all operations, except \mathcal{O}_1 and \mathcal{O}_2 , can be viewed as base operations which build the sparse skeleton of TDP-graphs, while \mathcal{O}_1 and \mathcal{O}_2 fill this skeleton with additional edges. This is also the main idea of the proof. First to discard all edges which are there by one of the operations \mathcal{O}_1 and \mathcal{O}_2 , and then study the resulting vertices of degree two.

Lemma 2.1. *If $(G, S) \in \mathcal{G}$ for some 2-coloring $S = (S_R, S_B)$, then G is a TDP-graph. Further, $S = (S_R, S_B)$ is a partition of $V(G)$ into two total dominating sets of G .*

Proof. We proceed by induction on the number, $k \geq 0$, of operations \mathcal{O}_1 through \mathcal{O}_{17} used to construct a 2-colored graph $(G, S) \in \mathcal{G}$. If $k = 0$, then (G, S) is one of the four 2-colored base graphs illustrated in Figure 1, and one can readily observe that G is a TDP-graph and $S = (S_R, S_B)$ is a partition of $V(G)$ into two total dominating sets of G . This establishes the base case. Let $k \geq 1$ and suppose that every 2-colored graph $(G', S') \in \mathcal{G}$ that can be constructed using fewer than k operations satisfies the desired result.

Let $(G, S) \in \mathcal{G}$ be a 2-colored graph that can be built from one of the 2-colored base graphs by a sequence of k operations $\mathcal{O}_1 - \mathcal{O}_{17}$. Let \mathcal{O}_j be the last operation of that sequence where $j \in [17]$, and let (G', S') be the graph obtained from the same 2-colored base graph with the same sequence as that used to construct (G, S) but without applying the last operation \mathcal{O}_j . Thus, $(G', S') \in \mathcal{G}$ can be constructed using fewer than k operations.

By the induction hypothesis, the graph G' is a TDP-graph and $S' = (S'_A, S'_B)$ is a partition of $V(G')$ into two total dominating sets of G' . If $j \in [2]$, then $S = S'$ and G is a TDP-graph since no new vertices were added. For $3 \leq j \leq 17$ it is a simple exercise to check from the color of the new vertices added to (G', S') when forming (G, S) that the operation \mathcal{O}_j yields two disjoint total domination sets, namely S_R and S_B . Thus, G is a TDP-graph, and $S = (S_R, S_B)$ is a partition of $V(G)$ into two total dominating sets of G . \square

3 Main result

Our main result is to provide a constructive characterization of the graphs that have two disjoint total dominating sets, or, equivalently, a characterization of the TDP-graphs.

We prove that the class of all TDP-graphs is precisely the family \mathcal{G} constructed in Section 2. A proof of Theorem 3.1 is given in Section 4.

Theorem 3.1. *A graph G is a TDP-graph if and only if every component of (G, S) is in \mathcal{G} for some 2-coloring S . Further, if $(G, S) \in \mathcal{G}$, then $S = (S_R, S_B)$ is a partition of $V(G)$ into two total dominating sets of G .*

4 Proof of Theorem 3.1

The sufficiency follows from Lemma 2.1. To prove the necessity, let G be a TDP-graph and let $S = (S_R, S_B)$ be a partition of $V(G)$ into two total dominating sets of G . We show that $(G, S) \in \mathcal{G}$ by induction on $m = |E(G)|$. Since G is a TDP-graph, we note that $\delta(G) \geq 2$, G has order $n \geq 4$, and $m \geq 4$. If $m = 4$, then necessarily $G \cong C_4$, and (G, S) is the 2-colored base graph G_1 , and so $(G, S) \in \mathcal{G}$. This establishes the base case. Let $m \geq 5$ and assume that every TDP-graph G' of size less than m where $S' = (S'_R, S'_B)$ is a partition of $V(G')$ into two total dominating sets satisfies $(G', S') \in \mathcal{G}$.

Let G be a TDP-graph of order n and size m , and let $S = (S_R, S_B)$ be a partition of $V(G)$ into two total dominating sets of G . If G is disconnected, we apply the inductive hypothesis to each component of G to produce the desired result. Hence, we may assume that G is connected.

Our general strategy in what follows is to reduce the graph G to a TDP-graph G' of size less than m , apply the inductive hypothesis to G' to show that $(G', S') \in \mathcal{G}$, and then reconstruct the graph (G, S) from (G', S') by applying one of the operations \mathcal{O}_x , $x \in [17]$, to show that $(G, S) \in \mathcal{G}$. We state this formally, since we will frequently use the following statement.

Statement 4.1. *If G' is a TDP-graph of size less than m , where $S' = (S'_R, S'_B)$ is a partition of $V(G')$ into two total dominating sets, and (G, S) can be constructed from (G', S') by applying one of the operations \mathcal{O}_x , where $x \in [17]$, then $(G, S) \in \mathcal{G}$.*

We define three graphs G_R , G_B and G_{RB} associated with the graph G and the partition $S = (S_R, S_B)$. Let G_R and G_B be the subgraphs of G induced by the sets S_R and S_B , respectively, and so $G_R = G[S_R]$ and $G_B = G[S_B]$. Let G_{RB} be the (spanning) subgraph of G with $V(G_{RB}) = V(G)$ and $E(G_{RB}) = E(G) \setminus (E(G_R) \cup E(G_B))$.

Claim 4.2. *If some component of G_R , G_B or G_{RB} is not a star, then $(G, S) \in \mathcal{G}$.*

Proof. Suppose that there exists a component, C , of G_R , G_B or G_{RB} which is not a star. If C contains a cycle $v_1 \dots v_k v_1$, $k \geq 3$, then G can be obtained from $G' = G - v_1 v_2$ by either applying operation \mathcal{O}_1 in the case when C is a component of G_R or G_B or by applying operation \mathcal{O}_2 in the case when C is a component of G_{RB} . If C contains no cycle, then C is a tree different from a star. Therefore, there exists a path $u_1 u_2 u_3 u_4$ in C and G can be obtained from $G' = G - u_2 u_3$ by either applying operation \mathcal{O}_1 in the case when C is a component of G_R or G_B or by applying operation \mathcal{O}_2 in the case when C is a component of G_{RB} . In all cases, since $S = (S_R, S_B)$ is a partition of $V(G)$ into two total dominating sets of G , the same partition $S' = S = (S_R, S_B)$ is a partition of $V(G')$ into two total

dominating sets of G' . By the inductive hypothesis, $(G', S') \in \mathcal{G}$. We can obtain G from the same 2-colored base graph as G' and the same sequence of operations from $\mathcal{O}_1 - \mathcal{O}_{17}$ used to construct (G', S') by adding at the end of this sequence the operation \mathcal{O}_1 or \mathcal{O}_2 . Hence $(G, S) \in \mathcal{G}$. \square

By Claim 4.2, we may assume that every component of G_R , G_B or G_{RB} is a star, for otherwise the desired result follows. We call the resulting graph G a *sparse TDP-graph with associated partition* $S = (S_R, S_B)$.

We now partition the sets S_R and S_B in two different ways depending on the role that the vertices in S_R and S_B , respectively, play in the graphs G_R , G_B and G_{RB} . First, let $S_R = R_1 \cup R_2 \cup R_3$ and $S_B = B_1 \cup B_2 \cup B_3$ where

$$\begin{aligned} R_1 &= \{v \in S_R \mid d_{G_R}(v) \geq 2\} \\ R_2 &= \{v \in S_R \setminus R_1 \mid N_G(v) \cap R_1 \neq \emptyset\} \\ R_3 &= S_R \setminus (R_1 \cup R_2) \end{aligned}$$

and

$$\begin{aligned} B_1 &= \{v \in S_B \mid d_{G_B}(v) \geq 2\} \\ B_2 &= \{v \in S_B \setminus B_1 \mid N_G(v) \cap B_1 \neq \emptyset\} \\ B_3 &= S_B \setminus (B_1 \cup B_2). \end{aligned}$$

Next, we define a partition of $V(G) = V(G_{RB})$ as the union of the two partitions $S_R = R_1B \cup R_2B \cup R_3B$ and $S_B = RB_1 \cup RB_2 \cup RB_3$ where

$$\begin{aligned} R_1B &= \{v \in S_R \mid d_{G_{RB}}(v) \geq 2\} \\ R_2B &= \{v \in S_R \setminus R_1B \mid v \text{ has a neighbor in } G_{RB} \text{ that belongs to } RB_1\} \\ R_3B &= S_R \setminus (R_1B \cup R_2B) \end{aligned}$$

and

$$\begin{aligned} RB_1 &= \{v \in S_B \mid d_{G_{RB}}(v) \geq 2\} \\ RB_2 &= \{v \in S_B \setminus RB_1 \mid v \text{ has a neighbor in } G_{RB} \text{ that belongs to } R_1B\} \\ RB_3 &= S_B \setminus (RB_1 \cup RB_2). \end{aligned}$$

We note that every vertex in R_3 has degree 1 in G_R , and every vertex in R_3B has degree 1 in G_{RB} . Analogously, every vertex in B_3 and RB_3 has degree 1 in G_B and G_{RB} , respectively. In particular, vertices from $R_3 \cap R_3B$ and from $B_3 \cap RB_3$ have degree 2 in G . Further, the neighbor of a vertex from R_3 in G_R belongs to R_3 , and, analogously, the neighbor of a vertex from B_3 in G_B belongs to B_3 . We proceed further with the following series of structural properties of the graph G .

Claim 4.3. $\delta(G) = 2$.

Proof. Recall that G is a sparse TDP-graph with associated partition $S = (S_R, S_B)$. Thus, S_R and S_B are disjoint total dominating sets of G which form a partition of $V(G)$. Every vertex $v \in V(G)$ has at least one neighbor in S_R and at least one neighbor in S_B . Hence, $\delta(G) \geq 2$. Suppose, to the contrary, that $\delta(G) > 2$.

Suppose that $R_1B \neq \emptyset$ and let $v \in R_1B$. Let v_1, \dots, v_k , where $k \geq 2$, be the neighbors of v in G_{RB} . By Claim 4.2 and the definition of the set RB_2 , we note that for each $i \in [k]$, $v_i \in RB_2$ and the vertex v is the only neighbor of v_i that belongs to the set S_R . Further, since $d_G(v_i) > 2$, the vertex v_i has at least two neighbors in S_B . By Claim 4.2, every component of the graph G_B is a star, implying that no two neighbors of v are adjacent or have a common neighbor in G_B . Further, every neighbor of v_i in G different from v belongs to the set B_2 , and has the vertex v_i as its only neighbor in G_B . Thus, the set B_2 contains at least $2k$ vertices at distance 2 from v in G .

For $i \in [k]$, let w_i denote an arbitrary neighbor of v_i in G_B , and so $w_i \in B_2$. Since $d_G(w_i) > 2$ and w_i has only one neighbor in S_B , namely the vertex v_i , we note that $w_i \in RB_1$ and therefore w_i has at least two neighbors in R_2B . By Claim 4.2 and the definition of the set R_2B , we note that every neighbor of w_i different from v_i belongs to the set R_2B . Further, each such neighbor of w_i has exactly one neighbor that belongs to the set S_B , namely the vertex w_i , and therefore has at least two neighbors in S_R by the minimum degree condition. By Claim 4.2, every component of the graph G_R is a star, and therefore two distinct vertices of degree at least 2 in G_R belong to different components of G_R . This implies that this subset R_2B of vertices in S_R contains at least $4k$ vertices.

By the minimum degree condition, these vertices in R_2B also belong to R_1 and each of them has at least two neighbors in R_2 . Further, analogously as before, no two such vertices are the same, implying that this subset of R_2 contains at least $8k - 1$ vertices distinct from v , all of which belong to the set R_1B , noting that one of these vertices may possibly be the vertex v . By repeating this process for all these vertices we see that we have an infinite process with infinite growth, which is not possible in a finite graph G . Therefore, the set $R_1B = \emptyset$. Analogously, the set $RB_1 = \emptyset$. Therefore, R_2B and RB_2 are also empty.

We now consider a vertex $v \in R_3B$. By Claim 4.2, every component of the graph G_{RB} is a star, implying that the vertex v has exactly one neighbor in S_B and, by the minimum degree condition, at least two neighbors in S_R . Thus, $v \in R_1$ and each neighbor of v in S_R belong to R_2 . Further, by Claim 4.2, each such neighbor of v in R_2 has degree 1 in G_R and, therefore, by the minimum degree condition, has at least two neighbors in S_B . Thus, every neighbor of v in R_2 belongs to the set R_1B , contradicting our earlier observation that the set R_1B is an empty set. This completes the proof of Claim 4.3. \square

By Claim 4.3, every sparse TDP-graph has minimum degree 2. In particular, $\delta(G) = 2$. Let $D = \{v \in V(G) \mid d_G(v) = 2\}$.

Claim 4.4. *If a vertex in D is a cut-vertex of G , then $(G, S) \in \mathcal{G}$.*

Proof. Suppose that a vertex in D is a cut-vertex of G . Suppose firstly that D contains two adjacent vertices, x and y , that are both cut-vertices of G , and let $e = xy$. Let C_x and C_y be the components of $G - e$ which contain x and y , respectively. Further, let x' be the neighbor of x in C_x and let y' be the neighbor of y in C_y . We have two possibilities with respect to the color of the vertices x, x', y, y' . Either $\text{color}(x') = \text{color}(x) \neq \text{color}(y) = \text{color}(y')$ or $\text{color}(x') = \text{color}(y') \neq \text{color}(x) = \text{color}(y)$. In both cases, let G' be the graph obtained from $G - \{x, y\}$ by adding the edge $x'y'$, and changing the color of all vertices in $V(C_y) \setminus \{y\}$ while retaining the color of all vertices in $V(C_x) \setminus \{x\}$. Let $S' = (S'_R, S'_B)$ be the resulting partition of $V(G')$. We note that G' is a TDP-graph, where $S' = (S'_R, S'_B)$ is a partition of $V(G')$ into two total dominating sets and that x' and y' are cut vertices of G' . If x and x' have the same color in G , then we use Statement 4.1 with the operation \mathcal{O}_{17}

and the cut vertex y' to show that $(G, S) \in \mathcal{G}$, while if x and x' have different color in G , we use Statement 4.1 with the operation \mathcal{O}_{17} and the cut vertex x' .

Thus, we may assume that no two adjacent vertices of D are both cut-vertices of G . Let v be a cut-vertex of G that belongs to D with neighbors u' and u'' . Without loss of generality we may assume that $\text{color}(v) = \text{color}(u'') \neq \text{color}(u')$. Let $C_{u'}$ and $C_{u''}$ be the components of $G - v$ containing u' and u'' , respectively. Since $S = (S_R, S_B)$ is a partition of $V(G)$ into two total dominating sets of G , there exists a neighbor of u' in $C_{u'}$ of the same color as u' and a neighbor of u'' in $C_{u''}$ whose color is different from that of u'' . Let G' be the graph obtained from $G - v$ by identifying the vertices u' and u'' into one new vertex u , and joining u to every neighbor of u' and u'' . Further, we assign to u the same color as that of u' , while we change the color of all vertices in $V(C_{u''}) \setminus \{u''\}$ and retain the color of all vertices in $V(C_{u'}) \setminus \{u'\}$. Let $S' = (S'_R, S'_B)$ be the resulting partition of $V(G')$. We note that G' is a TDP-graph, where $S' = (S'_R, S'_B)$ is a partition of $V(G')$ into two total dominating sets. We now use Statement 4.1 with the operation \mathcal{O}_{17} to show that $(G, S) \in \mathcal{G}$, where $H_1^u = C_{u'}$ and $H_2^u = C_{u''}$. \square

By Claim 4.4, we may assume that no vertex in D is a cut-vertex of G , for otherwise the desired result follows. We note that every vertex in D has one neighbor in S_R and one neighbor in S_B . Further, every component in $G[D]$ is a path or a cycle.

Claim 4.5. *Let C be a component of $G[D]$. If C is a cycle or if C is a path of order at least 5 or if C is a path of order 4 and the ends of C do not have a common neighbor, then $(G, S) \in \mathcal{G}$.*

Proof. Suppose that C is a cycle. Since G is a connected TDP-graph, this implies that $G \cong C_n$ where $n \equiv 0 \pmod{4}$. In this case, G can be obtained from the 2-colored base graph G_1 by repeated applications of operation \mathcal{O}_9 (or operation \mathcal{O}_{10}). Hence, we may assume that C is a path, for otherwise the desired result follows. Let C be the path $x_1 \dots x_k$, where $k \geq 4$. Let u be the neighbor of x_1 not on C . If $k \geq 5$, let $v = x_5$, while if $k = 4$, let v be the neighbor of x_4 not on C . By assumption, $u \neq v$. Let $X = \{x_1, x_2, x_3, x_4\}$.

Suppose first that $\text{color}(u) = \text{color}(x_1)$, implying that $\text{color}(x_2) = \text{color}(x_3) \neq \text{color}(x_4) = \text{color}(v) = \text{color}(x_1)$. If u and v are adjacent in G , let $G' = G - X$. In this case, the graph G' is a TDP-graph and we use Statement 4.1 with the operation \mathcal{O}_7 to show that $(G, S) \in \mathcal{G}$. If u and v are not adjacent in G , let G' be obtained from $G - X$ by adding the edge uv . Once again, the graph G' is a TDP-graph. We use Statement 4.1 with the operation \mathcal{O}_{10} to show that $(G, S) \in \mathcal{G}$.

Suppose next that $\text{color}(u) \neq \text{color}(x_1)$, implying that $\text{color}(x_2) = \text{color}(v) \neq \text{color}(x_3) = \text{color}(x_4) = \text{color}(u)$. If u and v are adjacent in G , let $G' = G - X$. In this case, the graph G' is a TDP-graph and we use Statement 4.1 with the operation \mathcal{O}_8 to show that $(G, S) \in \mathcal{G}$. If u and v are not adjacent in G , let G' be obtained from $G - X$ by adding the edge uv . Once again, the graph G' is a TDP-graph. We use Statement 4.1 with the operation \mathcal{O}_9 to show that $(G, S) \in \mathcal{G}$. \square

By Claim 4.5, we may assume that every component of $G[D]$ is a path-component of order at most 4, and that the ends of a path-component of $G[D]$ of order 4 have a common neighbor in G . In what follows we adopt the following notation. Let P be a path-component of $G[D]$, and so $P \cong P_k$ for some $k \in [4]$. Let P be the path $x_1 \dots x_k$, and let u and v be the vertices in G that do not belong to P and are adjacent to x_1 and x_k ,

respectively. We call u and v the vertices in $G - V(P)$ associated with the path P . By assumption, if $k = 4$, then $u = v$. We note that if $k = 1$, then $u \neq v$. We define next a good path-component.

Definition 4.6. A path-component P of $G[D]$ is a *good path-component* if $P \cong P_k$ where $k \in [3]$, and both u and v have neighbors of both colors in the graph $G^- = G - V(P)$, where u and v are the vertices in G^- associated with P .

Claim 4.7. *If $G[D]$ contains a good path-component, then $(G, S) \in \mathcal{G}$.*

Proof. Suppose that $G[D]$ contains a good path-component, $P: x_1 \dots x_k$. By definition, $k \in [3]$. Suppose that $k = 1$. Since P is a good path-component, the graph $G' = G - x_1$ is a TDP-graph. Furthermore, $\text{color}(u) \neq \text{color}(v)$ since G is a TDP-graph. We now use Statement 4.1 with the operation \mathcal{O}_3 to show that $(G, S) \in \mathcal{G}$.

Suppose that $k = 2$. Suppose that $\text{color}(x_1) = \text{color}(x_2)$. Then, $\text{color}(u) \neq \text{color}(x_1)$ and either $u = v$ or $u \neq v$ and $\text{color}(u) = \text{color}(v)$. In both cases, since P is a good path-component, the graph $G' = G - V(P)$ is a TDP-graph. If $u = v$, we use Statement 4.1 with the operation \mathcal{O}_{11} to show that $(G, S) \in \mathcal{G}$, while if $u \neq v$, we use Statement 4.1 with the operation \mathcal{O}_4 to show that $(G, S) \in \mathcal{G}$. Suppose that $\text{color}(x_1) \neq \text{color}(x_2)$. Then, $\text{color}(u) = \text{color}(x_1)$ and $\text{color}(v) = \text{color}(x_2)$. Since P is a good path-component, the graph $G' = G - V(P)$ is a TDP-graph, and we use Statement 4.1 with the operation \mathcal{O}_5 to show that $(G, S) \in \mathcal{G}$.

Suppose that $k = 3$. Without loss of generality we may assume that $\text{color}(x_1) \neq \text{color}(x_2) = \text{color}(x_3)$, implying that $\text{color}(u) = \text{color}(x_1)$ and either $u = v$ or $u \neq v$ and $\text{color}(u) = \text{color}(v)$. Since P is a good path-component, the graph $G' = G - V(P)$ is a TDP-graph. If $u = v$, we use Statement 4.1 with the operation \mathcal{O}_{12} to show that $(G, S) \in \mathcal{G}$, while if $u \neq v$, we use Statement 4.1 with the operation \mathcal{O}_6 to show that $(G, S) \in \mathcal{G}$. □

By Claim 4.7, we may assume that G contains no good path-component, for otherwise the desired result follows. We define next an end-block path component of $G[D]$.

Definition 4.8. A path-component P of $G[D]$ with associated vertices u and v is an *end-block path component* of $G[D]$ if $u = v$.

We are now in a position to present the following property of non-backtracking walks in the graph G .

Claim 4.9. *Suppose that $W: w_1 w_2 \dots w_k$ is a non-backtracking walk in G and no vertex of W belongs to an end-block path component of $G[D]$. If w_2 is not the only neighbor of w_1 in G whose color is $\text{color}(w_2)$, then w_{i-1} is the only neighbor of w_i in G whose color is $\text{color}(w_{i-1})$ for all $i \in [k] \setminus \{1\}$.*

Proof. Since W is a non-backtracking walk in G , we note that no two consecutive edges on W are equal; that is, $w_{i-1} \neq w_{i+1}$ for all $i \in [k-1] \setminus \{1\}$. Suppose, to the contrary, that the claim is false. Let $\ell \geq 2$ be the smallest integer such that the vertex w_ℓ has a neighbor different from $w_{\ell-1}$ of the same color as $w_{\ell-1}$.

Claim 4.9.1. $\ell \geq 3$.

Proof. Renaming colors if necessary, we may assume that $\text{color}(w_1) = X$. By supposition, at least one neighbor, say v_1 , of w_1 different from w_2 has the same color as w_2 . Suppose firstly that $\text{color}(w_2) = X$. By supposition, $\text{color}(v_1) = X$. If w_2 has a neighbor, z_2 say, different from w_1 , of color X , then either $v_1 = z_2$, in which case $v_1 w_1 w_2 v_1$ is a 3-cycle in G_X , or $v_1 \neq z_2$, in which case $v_1 w_1 w_2 z_2$ is a path P_4 in G_X . Both cases produce a contradiction. Suppose secondly that $\text{color}(w_2) = \bar{X}$. By supposition, $\text{color}(v_1) = \bar{X}$. If w_2 has a neighbor, z_2 say, different from w_1 , of color X , then $v_1 w_1 w_2 z_2$ is a path P_4 in G_{RB} , a contradiction. We deduce, therefore, that w_1 is the only neighbor of w_2 whose color is $\text{color}(w_1)$. Hence, $\ell \geq 3$. \square

By Claim 4.9.1, we have that $\ell \geq 3$. Renaming colors if necessary, we may assume that $\text{color}(w_{\ell-1}) = X$. By supposition, the vertex w_ℓ has a neighbor, $v_{\ell+1}$ say, different from $w_{\ell-1}$ of the same color as $w_{\ell-1}$; that is, $\text{color}(v_{\ell+1}) = X$. Further since G is a TPD-graph, the vertex w_ℓ has a neighbor of color \bar{X} .

Claim 4.9.2. $d_G(w_{\ell-1}) = 2$.

Proof. Suppose that $d_G(w_{\ell-1}) \geq 3$. Let v_ℓ be a neighbor of $w_{\ell-1}$ different from $w_{\ell-2}$ and w_ℓ . Suppose that $\text{color}(w_{\ell-2}) = X$. By the minimality of ℓ , the vertex $w_{\ell-2}$ is the only neighbor of $w_{\ell-1}$ whose color is $\text{color}(w_{\ell-2})$; that is, all neighbors of $w_{\ell-1}$ different from $w_{\ell-2}$ must have color \bar{X} . In particular, $\text{color}(v_\ell) = \text{color}(w_\ell) = \bar{X}$. Hence, $v_\ell w_{\ell-1} w_\ell v_{\ell+1}$ is a path P_4 in G_{RB} , a contradiction. Hence, $\text{color}(w_{\ell-2}) = \bar{X}$. Thus, all neighbors of $w_{\ell-1}$ different from $w_{\ell-2}$ must have color X . In particular, $\text{color}(v_\ell) = \text{color}(w_\ell) = X$. If $v_\ell = v_{\ell+1}$, then $v_\ell w_{\ell-1} w_\ell v_\ell$ is a 3-cycle in G_X , a contradiction. If $v_\ell \neq v_{\ell+1}$, then $v_\ell w_{\ell-1} w_\ell v_{\ell+1}$ is a path P_4 in G_X , a contradiction. \square

By Claim 4.9.2, the vertex $w_{\ell-1}$ has degree 2 in G ; that is, $w_{\ell-1} \in D$. By supposition, the vertex w_1 has at least two neighbors whose color is $\text{color}(w_2)$ and at least one vertex whose color is different from $\text{color}(w_2)$. In particular, the vertex w_1 has degree at least 3 in G . Let $p \geq 1$ be the largest integer such that $d_G(w_p) \geq 3$ and $p \leq \ell - 2$. Possibly, $p = \ell - 2$. We now consider the path $P : w_{p+1} \dots w_{\ell-1}$ and note that P is a path-component in $G[D]$. If $w_p = w_\ell$, then P is an end-block path component of $G[D]$, contradicting the supposition that no vertex of W belongs to an end-block path component of $G[D]$. Hence, $w_p \neq w_\ell$ and the vertices w_p and w_ℓ associated with the path-component P in $G[D]$ are distinct vertices.

We now consider the graph $G^- = G - V(P)$. By our earlier observations, the vertex w_ℓ has neighbors of both colors in G^- . If $p = 1$, then the vertex w_p has neighbors of both colors in G^- . If $p \geq 2$, then by the minimality of ℓ the vertex w_p once again has neighbors of both colors in G^- . Thus the path P is a good-path component, contradicting our earlier assumption that G contains no good path-component. This completes the proof of Claim 4.9. \square

Claim 4.10. *If G contains a cycle that is not an end-block of G , then $(G, S) \in \mathcal{G}$.*

Proof. Assume that some cycle C in G is not an end-block in G . Let P be a path-component of $G[D]$ with associated vertices u and v . Suppose firstly that $u = v$. Thus, P is an end-block path component of $G[D]$ and $C_P = G[V(P) \cup \{u\}]$ is a cycle in G . Further, C_P is an end-block of G with u as its cut-vertex in G . Suppose that $d_G(u) \geq 4$. We now consider the graph $G^- = G - V(P)$. By our earlier assumptions, no vertex in D is a cut-vertex of G , implying that G^- is a connected graph.

Claim 4.10.1. *The vertex u has neighbors of both colors in G^- .*

Proof. Suppose, to the contrary, that all neighbors of u in G^- have the same color. By supposition, there is a cycle C in G^- that contains no vertex that belongs to an end-block component of $G[D]$. Hence there exists a non-backtracking walk $W: w_1w_2 \dots w_k$ in G that starts at the vertex u , proceeds from u to C , goes all the way around C , and then returns to u , without entering any end-block path component of $G[D]$. We note that $k \geq 3$ and that $w_1 = w_k = u$. By our supposition that all neighbors of u in G^- have the same color, the vertex w_2 is not the only neighbor of w_1 in G whose color is $\text{color}(w_2)$. By Claim 4.9, the vertex w_{k-1} is the only neighbor of w_k in G whose color is $\text{color}(w_{k-1})$. This contradicts our supposition that all neighbors of u in G^- have the same color. \square

By Claim 4.10.1, the vertex u has neighbors of both colors in G^- . Since G is a TDP-graph, this implies that the graph G^- is a TDP-graph. Hence, we can use Statement 4.1 with the operation \mathcal{O}_{11} or \mathcal{O}_{12} or \mathcal{O}_{13} , depending on the length of P , to show that $(G, S) \in \mathcal{G}$. We may therefore assume that $d_G(u) = 3$ (and still $u = v$), for otherwise $(G, S) \in \mathcal{G}$, as desired. Thus, the vertex u has degree 1 in G^- . Let x be the neighbor of u in G^- . By our earlier assumptions, no vertex in D is a cut-vertex of G . In particular, the cut-vertex x does not belong to D , and so $d_G(x) \geq 3$. We now consider the (connected) graph $G_u^- = G^- - u$ obtained from G^- by deleting the vertex u . Using analogous arguments as in the proof of Claim 4.10.1, the vertex x has neighbors of both colors in G_u^- . Hence, we can use Statement 4.1 with the operation \mathcal{O}_{14} or \mathcal{O}_{15} or \mathcal{O}_{16} , depending on the length of P , to show that $(G, S) \in \mathcal{G}$.

Suppose next that $u \neq v$. Using analogous arguments as in the proof of Claim 4.10.1, the vertices u and v each have neighbors of both colors in G_u^- . Thus the path P is a good-path component, contradicting our earlier assumption that G contains no good path-component. This completes the proof of Claim 4.10. \square

By Claim 4.10, we may assume that every cycle in G is an end-block of G , for otherwise $(G, S) \in \mathcal{G}$ as desired. Every block of G that is not an end-block is a copy of K_2 consisting of a single edge. By our earlier assumptions, every cycle in G has length 3, 4 or 5. Let T^- be the graph obtained from G by deleting all vertices that belong to an end-block path component of $G[D]$. By our earlier assumptions, the graph T^- is a tree. In particular, every vertex of T^- is a cut-vertex of G . By our earlier assumptions, no vertex in D is a cut-vertex of G , implying that every vertex of D belongs to an end-block path component of $G[D]$. Hence, every vertex of D belongs to an end-block of G .

Claim 4.11. *If two cycles of G intersect, then $(G, S) \in \mathcal{G}$.*

Proof. Suppose that two different cycles C_1 and C_2 of G intersect. Since every cycle in G is an end-block of G , these two cycles intersect in exactly one common vertex, v say.

Claim 4.11.1. *If G contains exactly one cut-vertex, then (G, S) is a 2-colored base graph G_3 .*

Proof. Suppose that G contains exactly one cut-vertex. Since the cut-vertices of G are precisely the vertices in the tree T^- , this implies that $V(T^-) = \{v\}$. Thus, every block of G is an end-block that contains the vertex u . Let C_1 be the cycle $vv_1v_2 \dots v_{k-1}v$ and let $\text{color}(v) = X$, where $k \in \{3, 4, 5\}$. If $k = 3$, then $\text{color}(v_1) = \text{color}(v_2) = \overline{X}$. If $k = 4$, then $\text{color}(v_2) = \overline{X}$ and, renaming v_1 and v_3 , if necessary, we may assume that

$\text{color}(v_1) = X$ and $\text{color}(v_3) = \bar{X}$. If $k = 5$, then $\text{color}(v_2) = \text{color}(v_3) = \bar{X}$ and $\text{color}(v_1) = \text{color}(v_4) = X$.

Suppose that G contains an end-block, C say, that is a 4-cycle. If C' is an arbitrary end-block different from C , then $C' - v$ is a good path-component of $G[D]$, a contradiction. Hence, no end-block of G is a 4-cycle.

Thus, since G is a TDP-graph, at least one end-block is a 3-cycle and at least one end-block is a 5-cycle. Renaming the end-blocks if necessary, we may assume that C_1 is a 3-cycle and C_2 is a 5-cycle. These two cycles, together with their associated 2-colorings described above, form the 2-colored base graph G_3 . If G contains at least three blocks and C' is an arbitrary end-block different from C_1 and C_2 , then $C' - v$ is a good path-component of $G[D]$, a contradiction. Hence, G contains exactly two end-blocks, implying that (G, S) is the 2-colored base graph G_3 . \square

By Claim 4.11.1, we may assume that G contains at least two cut-vertices, for otherwise $(G, S) \in \mathcal{G}$ as desired. As observed earlier, the cut-vertices of G are precisely the vertices in the tree T^- . Let x be a neighbor of v in T^- . Renaming the cycle C_1 and C_2 and the vertex x if necessary, we may assume without loss of generality that the vertex v has a neighbor, y say, in C_1 such that $\text{color}(x) \neq \text{color}(y)$. We now consider the graph $G^- = G - (V(C_2) \setminus \{v\})$. Since G is a TDP-graph, this implies that the graph G^- is a TDP-graph. Hence, we can use Statement 4.1 with the operation \mathcal{O}_{11} or \mathcal{O}_{12} or \mathcal{O}_{13} , depending on the length of C_2 , to show that $(G, S) \in \mathcal{G}$. \square

By Claim 4.11, we may assume that no two cycles of G intersect, for otherwise $(G, S) \in \mathcal{G}$ as desired. The tree T^- therefore contains at least two vertices. Further, every leaf in T^- has degree 3 in G and belongs to exactly one end-block of G . Let $p_1 p_2 \dots p_k$ be a longest path in T^- . Necessarily, p_1 and p_k are both leaves in T^- . Since T^- contains no vertex of D , we note that every vertex in T^- has degree at least 3 in G . Let C_1 and C_k be the end-blocks in G that contain p_1 and p_k , respectively.

Claim 4.12. *If $k \in \{2, 3\}$, then $(G, S) \in \mathcal{G}$.*

Proof. Suppose firstly that $k = 2$. In this case, G is obtained from the two cycles C_1 and C_2 by adding the edge $p_1 p_2$. If C_1 is a 4-cycle, then the cycle C_1 together with its associated 2-coloring is the 2-colored base graph G_1 . Starting with this 2-colored base graph G_1 , we can use Statement 4.1 with the operation \mathcal{O}_{14} or \mathcal{O}_{15} or \mathcal{O}_{16} , depending on the length of C_2 , to show that $(G, S) \in \mathcal{G}$. Analogously, if C_2 is a 4-cycle, $(G, S) \in \mathcal{G}$. Hence, we may assume that neither C_1 nor C_2 is a 4-cycle. With this assumption, if C_1 is a 3-cycle, then C_2 is also a 3-cycle noting that G is a TDP-graph. In this case, (G, S) is the 2-colored base graph G_2 . If C_1 is a 5-cycle, then C_2 is also a 5-cycle. In this case, (G, S) is the 2-colored base graph G_4 . Hence if $k = 2$, then $(G, S) \in \mathcal{G}$.

Suppose secondly that $k = 3$. We now consider the (connected) graph $G^- = G - V(C_1)$. We note that the vertex p_2 has degree at least 2 in G^- . If the vertex p_2 has neighbors of both colors in G^- , then G^- is a TPD-graph. In this case, we can use Statement 4.1 with the operation \mathcal{O}_{14} or \mathcal{O}_{15} or \mathcal{O}_{16} , depending on the length of C_1 , to show that $(G, S) \in \mathcal{G}$. Hence we may assume that all neighbors of p_2 in G^- have the same color which is different to $\text{color}(p_1)$ (noting that G is a TPD-graph). This implies that the vertex p_2 has neighbors of both colors in the graph $G - V(C_2)$, and once again we can use Statement 4.1 with the operation \mathcal{O}_{14} or \mathcal{O}_{15} or \mathcal{O}_{16} , depending on the length of C_2 , to show that $(G, S) \in \mathcal{G}$. \square

By Claim 4.12, we may assume that $k \geq 4$, for otherwise $(G, S) \in \mathcal{G}$ as desired. We now consider the (connected) graph $G^- = G - V(C_1)$. If the vertex p_2 has neighbors of both colors in G^- , then as in the proof of Claim 4.12 we can use Statement 4.1 with the operation \mathcal{O}_{14} or \mathcal{O}_{15} or \mathcal{O}_{16} , depending on the length of C_1 , to show that $(G, S) \in \mathcal{G}$. Hence we may assume that all neighbors of p_2 in G^- have the same color. We now consider the walk $p_2 p_3 \dots p_k$. By assumption, p_3 is not the only neighbor of p_2 in G whose color is $\text{color}(p_3)$. By Claim 4.9, the vertex p_{k-2} is the only neighbor of p_{k-1} in G whose color is $\text{color}(p_{k-2})$. This implies that the vertex p_{k-1} has neighbors of both colors in the graph $G - V(C_2)$. Hence, $G - V(C_2)$ is a TPD-graph and we can use Statement 4.1 with the operation \mathcal{O}_{14} or \mathcal{O}_{15} or \mathcal{O}_{16} , depending on the length of C_2 , to show that $(G, S) \in \mathcal{G}$. This completes the proof of Theorem 3.1. \square

5 Closing remarks

We remark that although our characterization in Theorem 3.1 solves a long-standing problem in the theory of total domination in graphs which has been open for several decades, it remains a challenging problem to determine in polynomial time if a given graph is a TDP-graph even for some special graph classes. Our method cannot be used to decide if a given graph G is a TDP-graph in polynomial time. The reason for that is that we have no specified vertex partition together with G . Indeed, recognizing this class of graphs is known to be NP-complete (see [8]). However, we nonetheless believe that our constructive proof gives valuable insights into the problem and gives an entirely new description of TDP-graphs, placing them in another context.

We close with a short discussion about the independence of operations \mathcal{O}_1 to \mathcal{O}_{17} in the class \mathcal{G} . For this purpose, we will construct small graphs in \mathcal{G} from our 2-colored base graphs that cannot be built by any other construction in \mathcal{G} , thereby showing that operation \mathcal{O}_i is independent for each $i \in [17]$. The independence of these seventeen operations used to build graphs in the family \mathcal{G} show that none of them are redundant, and all are needed in the construction.

- Apply operation \mathcal{O}_2 on G_1 (to obtain the graph $K_4 - e$).
- Apply operation \mathcal{O}_3 on G_1 to obtain the house graph; that is, the graph obtained from a 5-cycle by adding an edge.
- Apply operation \mathcal{O}_1 once and operation \mathcal{O}_2 three times on the house graph to obtain K_5 .
- Apply operation \mathcal{O}_4 to two nonadjacent vertices of degree 2 on G_2 .
- The independence of operation \mathcal{O}_x , where $x \in \{5, 6, 11, 12, 13, 14, 15, 16\}$, can be seen by applying \mathcal{O}_x once on G_1 .
- The independence of operation \mathcal{O}_x , where $x \in \{7, 10\}$, can be seen by applying \mathcal{O}_x once on adjacent vertices of degree 3 in G_2 .
- The independence of operation \mathcal{O}_x , where $x \in \{8, 9\}$, can be seen by applying \mathcal{O}_x once on adjacent vertices of degree 3 in G_4 .
- Apply operation \mathcal{O}_{17} once on the cut-vertex of G_3 .

Hence, all seventeen operations are independent. Further, our proof of Theorem 3.1 shows that all seventeen operations are necessary to give our characterization of TDP-graphs.

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